
A UNIFIED APPROACH TO DISTANCE-TWO COLOURING OF GRAPHS ON SURFACES

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Abstract. — In this paper we introduce the notion of Σ -colouring of a graph G : For given subsets $\Sigma(v)$ of neighbours of v , for every $v \in V(G)$, this is a proper colouring of the vertices of G such that, in addition, vertices that appear together in some $\Sigma(v)$ receive different colours. This concept generalises the notion of colouring the square of graphs and of cyclic colouring of graphs embedded in a surface. We prove a general result for graphs embeddable in a fixed surface, which implies asymptotic versions of Wegner’s and Borodin’s Conjecture on the planar version of these two colourings. Using a recent approach of Havet *et al.*, we reduce the problem to edge-colouring of multigraphs, and then use Kahn’s result that the list chromatic index is close to the fractional chromatic index.

Our results are based on a strong structural lemma for graphs embeddable in a fixed surface, which also implies that the size of a clique in the square of a graph of maximum degree Δ embeddable in some fixed surface is at most $\frac{3}{2}\Delta$ plus a constant.

1. Introduction

Most of the terminology and notation we use in this paper is standard and can be found in any text book on graph theory (such as [2] or [8]). All our graphs and multigraphs will be finite. A *multigraph* can have multiple edges; a *graph* is supposed to be simple. We will not allow loops. The vertex and edge set of a graph G are denoted by $V(G)$ and $E(G)$, respectively (or just V and E , if the graph G is clear from the context).

Given a graph G , the *chromatic number* of G , denoted $\chi(G)$, is the minimum number of colours required so that we can properly colour its vertices using those colours. If we colour the edges of G , we get the *chromatic index*, denoted $\chi'(G)$. The *list chromatic number* or *choice number* $ch(G)$ is the minimum value k such that if we give each vertex v of G a list $L(v)$ of at least k colours, then we can find a proper colouring in which each vertex gets assigned a colour from its own private list. The list chromatic index $ch'(G)$ is defined analogously for edges.

The *square* G^2 of a graph G is the graph with vertex set $V(G)$, with an edge between any two different vertices that have distance at most two in G . A proper vertex colouring of the square of a graph can also be seen as a vertex colouring of the original graph satisfying:

- vertices that are adjacent receive different colours, and
- vertices that have a common neighbour receive different colours.

Another way to formulate these conditions is as ‘vertices at distance one or two must receive different colours’. This is why the name *distance-two colouring* is also used in the literature.

In this paper we consider a colouring concept that generalises the concept of colouring the square of a graph, but that also can be used to study different concepts such as *cyclic colouring of plane graphs* (definition will be given later).

For a vertex $v \in V$, let $N(v)$ (or $N_G(v)$ if we want to specify the graph under consideration) be the set of vertices adjacent to v . Suppose that for each vertex $v \in V$, we are given a subset $\Sigma(v) \subseteq N(v)$ of its neighbourhood. We call such a collection a Σ -system for G .

A Σ -colouring of G is an assignment of colours to the vertices of G so that:

- vertices that are adjacent receive different colours, and
- vertices that appear together in some $\Sigma(v)$ receive different colours.

When additionally each vertex v has its own list $L(v)$ of colours from which its colour must be chosen, we talk about a *list Σ -colouring*.

We denote by $\chi(G; \Sigma)$ the minimum number of colours required for a Σ -colouring to exist. Its list variant is denoted by $ch(G; \Sigma)$, and is defined as the minimum integer k such that for each assignment of a list $L(v)$ of at least k colours to vertices $v \in V$, there exists a proper Σ -colouring of G in which all vertices are assigned colours from their own lists.

Notice that we trivially have $\chi(G) = \chi(G; \emptyset)$ and $\chi(G^2) = \chi(G; N_G)$; and the same relations holds for the list variant (\emptyset assigns the empty set to each vertex).

We define the *width* of a Σ -system of G as $\Delta(G; \Sigma) = \max_{v \in V} |\Sigma(v)|$. It is clear that we always need at least $\Delta(G; \Sigma) + 1$ colours in a proper Σ -colouring. In the case $\Sigma \equiv N_G$, there exist plenty of graphs G that require $O(\Delta(G)^2)$ colours (where $\Delta(G) = \Delta(G; N_G)$ is the usual maximum degree of G). But for planar graphs, it is known that a constant times $\Delta(G)$ colours is enough (even for list colouring). We will take a closer look at this in Subsection 1.1 below.

Following Wegner’s Conjecture on colouring the square of planar graphs (see also next subsection), we propose the following conjecture.

Conjecture 1.1. — *There exist constants c_1, c_2 and c_3 such that for all planar graphs G and any Σ -system for G , we have*

$$\begin{aligned} \chi(G; \Sigma) &\leq \lfloor \frac{3}{2} \Delta(G; \Sigma) \rfloor + c_1; \\ ch(G; \Sigma) &\leq \lfloor \frac{3}{2} \Delta(G; \Sigma) \rfloor + c_2; \\ ch(G; \Sigma) &\leq \lfloor \frac{3}{2} \Delta(G; \Sigma) \rfloor + 1, \quad \text{if } \Delta(G; \Sigma) \geq c_3. \end{aligned}$$

If $\Sigma \equiv \emptyset$ (hence $\Delta(G; \Sigma) = 0$), then the Four Colour Theorem implies that the smallest possible value for c_1 is four; while the fact that planar graphs are always 5-list colourable but not always 4-list colourable, shows that the smallest possible value for c_2 is five.

Our main result is that Conjecture 1.1 is asymptotically correct: $ch(G; \Sigma) \leq \frac{3}{2} \Delta(G; \Sigma) + o(\Delta(G; \Sigma))$. In fact, we can prove this asymptotic result holds for general surfaces.

Theorem 1.2. — *For every surface S and real $\varepsilon > 0$, there exists a constant $\beta_{S, \varepsilon}$ such that the following holds for all $\beta \geq \beta_{S, \varepsilon}$. If G is a graph embeddable in S , with a Σ -system of width at most β , then $ch(G; \Sigma) \leq (\frac{3}{2} + \varepsilon) \beta$.*

A trivial lower bound for the (list) chromatic number of a graph G is the *clique number* $\omega(G)$, the maximum size of a clique in G . For graphs with a Σ -system, we can define the following related concept. A Σ -clique is a subset $C \subseteq V$ such that every two different vertices in C

are adjacent or appear together in some $\Sigma(v)$. Denote by $\omega(G; \Sigma)$ the maximum size of a Σ -clique in G . Then we trivially have $ch(G; \Sigma) \geq \omega(G; \Sigma)$, and so Theorem 1.2 means that for a graph G embeddable in some fixed surface S , we have $\omega(G; \Sigma) \leq \frac{3}{2} \Delta(G; \Sigma) + o(\Delta(G; \Sigma))$.

But in fact, the structural result we use to prove Theorem 1.2 fairly easily gives $\omega(G; \Sigma) \leq \frac{3}{2} \Delta(G; \Sigma) + O(1)$.

Theorem 1.3. — *For every surface S , there exist constants β_S and γ_S such that the following holds for all $\beta \geq \beta_S$. If G is a graph embeddable in S , with a Σ -system of width at most β , then every Σ -clique in G has size at most $\frac{3}{2}\beta + \gamma_S$.*

The main steps in the proof of Theorem 1.2 can be found in Section 2. The proof relies on two technical lemmas; the proofs of those can be found in Section 3. After that we use one of those lemmas to provide the relatively short proof of Theorem 1.3 in Section 4. In Section 5 we discuss some of the aspects of our work and discuss open problems related to (list) Σ -colouring of graphs. The final section provides some background regarding the proof by Kahn [17] of the asymptotical equality of the fractional chromatic index and the list chromatic index of multigraphs. A more general result, contained implicitly in Kahn’s work, is of crucial importance to our proof in this paper.

In the next two subsections, we discuss two special consequences of these results. These special versions of Theorems 1.2 and 1.3 also show that the term $\frac{3}{2}\beta$ is best possible.

But before presenting these applications, a remark is in order. In an earlier version of this paper, we gave our results in terms of (A, B) -colourings. For a graph G and vertex sets $A, B \subseteq V$ (not necessarily disjoint), an (A, B) -colouring of G is a colouring of the vertices in B such that adjacent vertices, and vertices with a common neighbour in A , receive different colours.

There is an obvious way to translate an (A, B) -colouring problem into a Σ -colouring problem: For $v \in A$ set $\Sigma(v) = N_G(v) \cap B$, and for $v \notin A$ set $\Sigma(v) = \emptyset$. Note that after this translation we are required to colour all vertices, not just those in B . But the vertices outside B do not appear in any $\Sigma(v)$, hence colouring them for a graph embeddable in a fixed surface requires at most a constant number of colours.

On the other hand, it is easy to construct instances of Σ -colouring problems for which there is no obvious translation to an (A, B) -colouring problem. In that sense, we feel justified in considering Σ -colouring as a more general concept. Moreover, the concept is general enough to allow a simplification of several arguments in Section 2, compared to the earlier version of our results.

1.1. Colouring the Square of Graphs. — Recall that the *square* of a graph G , denoted G^2 , is the graph with the same vertex set as G and with an edge between any two different vertices that have distance at most two in G . If G has maximum degree Δ , then a vertex colouring of its square will need at least $\Delta + 1$ colours, and the greedy algorithm shows that it is always possible to find a colouring of G^2 with $\Delta^2 + 1$ colours. Cages of diameter two, such as the 5-cycle, the Petersen graph and the Hoffman-Singleton graph (see, e.g., [2, page 84]), show that there exist graphs that in fact require $\Delta^2 + 1$ colours.

Regarding the chromatic number of the square of a planar graph, Wegner [33] posed the following conjecture (see also the book of Jensen and Toft [14, Section 2.18]), suggesting that for planar graphs far less than $\Delta^2 + 1$ colours suffice.

Conjecture 1.4 (Wegner [33]). — For a planar graph G of maximum degree Δ , $\chi(G^2) \leq$

$$\begin{cases} 7, & \text{if } \Delta = 3, \\ \Delta + 5, & \text{if } 4 \leq \Delta \leq 7, \\ \lfloor \frac{3}{2} \Delta \rfloor + 1, & \text{if } \Delta \geq 8. \end{cases}$$

Wegner also gave examples showing that these bounds would be tight. For even $\Delta \geq 8$, these examples are sketched in Figure 1(a). The graph in the picture has maximum degree $2k$ and

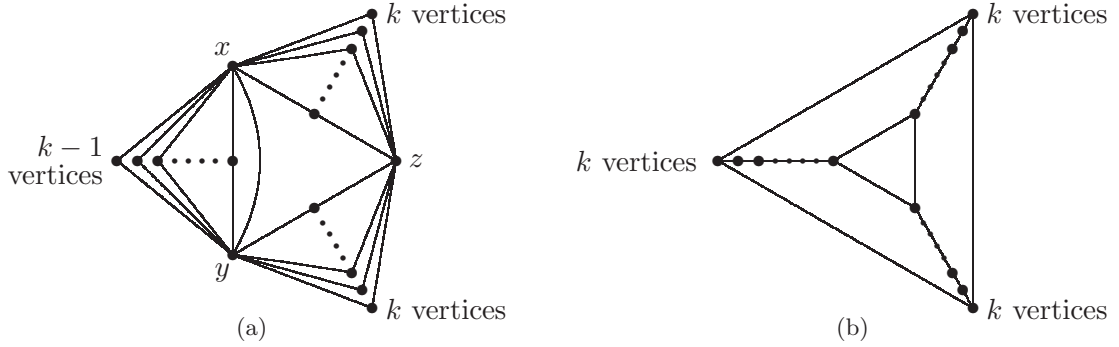


FIGURE 1. (a) A planar graph G with maximum degree $\Delta = 2k$ and $\omega(G^2) = \chi(G^2) = 3k + 1 = \lfloor \frac{3}{2} \Delta \rfloor + 1$.
 (b) A planar graph H with maximum face order $\Delta^* = 2k$ and $\chi^*(H) = 3k = \lfloor \frac{3}{2} \Delta^* \rfloor$ (see Subsection 1.2).

yet all the vertices except z are pairwise adjacent in its square. Hence to colour these $3k + 1$ vertices, we need at least $3k + 1 = \frac{3}{2} \Delta + 1$ colours. Note that the same arguments also show that the graph G in the picture has $\omega(G^2) = \frac{3}{2} \Delta + 1$.

Kostochka and Woodall [19] conjectured that for every square of a graph, the chromatic number equals the list chromatic number. This conjecture and Wegner's one together imply the conjecture that for planar graphs G with $\Delta \geq 8$, we have $ch(G^2) \leq \lfloor \frac{3}{2} \Delta \rfloor + 1$.

The first upper bound on $\chi(G^2)$ for planar graphs in terms of Δ , $\chi(G^2) \leq 8\Delta - 22$, was implicit in the work of Jonas [15]. This bound was later improved by Wong [34] to $\chi(G^2) \leq 3\Delta + 5$ and then by Van den Heuvel and McGuinness [13] to $\chi(G^2) \leq 2\Delta + 25$. Better bounds were then obtained for large values of Δ . It was shown that $\chi(G^2) \leq \lceil \frac{9}{5} \Delta \rceil + 1$ for $\Delta \geq 750$ by Agnarsson and Halldórsson [1], and the same bound for $\Delta \geq 47$ by Borodin *et al.* [4]. Finally, the best known upper bound so far has been obtained by Molloy and Salavatipour [25]: $\chi(G^2) \leq \lceil \frac{5}{3} \Delta \rceil + 78$. As mentioned in [25], the constant 78 can be reduced for sufficiently large Δ . For example, it was improved to 24 when $\Delta \geq 241$.

Since $ch(G^2) = ch(G; N_G)$ (i.e., $\Sigma(v) = N_G(v)$ for all $v \in V$), as an immediate corollary of Theorem 1.2 we obtain.

Corollary 1.5. — Let S be a fixed surface. Then the square of every graph G embeddable in S and of maximum degree Δ has list chromatic number at most $\frac{3}{2} \Delta + o(\Delta)$.

In fact, the same asymptotic upper bound as in Corollary 1.5 can be proved even for larger classes of graphs. Additionally, a stronger conclusion on the colouring is possible. For the following result, we assume that colours are integers, which allows us to talk about the 'distance' $|\alpha_1 - \alpha_2|$ between two colours α_1, α_2 .

Theorem 1.6 (Havet, Van den Heuvel, McDiarmid & Reed [10])

Let k be a fixed positive integer. The square of every $K_{3,k}$ -minor free graph G of maximum degree Δ has list chromatic number (and hence clique number) at most $\frac{3}{2}\Delta + o(\Delta)$. Moreover, given lists of this size, there is a proper colouring in which the colours on every pair of adjacent vertices of G differ by at least $\Delta^{1/4}$.

Note that planar graphs do not have a $K_{3,3}$ -minor. In fact, for every surface S , there is a constant k such that no graph embeddable in S has $K_{3,k}$ as a minor. That shows that Theorem 1.6 is stronger than our Corollary 1.5. On the other hand, Theorem 1.6 gives a weaker bound for the clique number than the one we obtain in Corollary 1.7 below.

Both Corollary 1.5 and Theorem 1.6 can be applied to K_4 -minor free graphs, since these graphs are planar and do not have $K_{3,3}$ as a minor. But the best possible bounds for this class are actually known. Lih, Wang and Zhu [21] showed that the square of K_4 -minor free graphs with maximum degree Δ has chromatic number at most $\lfloor \frac{3}{2}\Delta \rfloor + 1$ if $\Delta \geq 4$ and $\Delta + 3$ if $\Delta = 2, 3$. The same bounds, but then for the list chromatic number of K_4 -minor free graphs, were proved by Hetherington and Woodall [12].

Regarding the clique number of the square of graphs, we get the following corollary of Theorem 1.3.

Corollary 1.7. — *Let S be a fixed surface. Then the square of every graph G embeddable in S and of maximum degree Δ has clique number at most $\frac{3}{2}\Delta + O(1)$.*

From the proof of Theorem 1.3, it can be deduced that the square of a planar graph with maximum degree $\Delta \geq 11616$ has clique number at most $\frac{3}{2}\Delta + 76$.

Very recently, this was improved by the following result.

Theorem 1.8 (Cohen & Van den Heuvel [7]). — *For a planar graph G of maximum degree $\Delta \geq 41$, we have $\omega(G^2) \leq \lfloor \frac{3}{2}\Delta \rfloor + 1$.*

Apart from the bound $\Delta \geq 41$, this theorem is best possible, as is shown by the same graphs that show that Wegner’s Conjecture 1.4 is best possible for $\Delta \geq 8$ (see also Figure 1(a)).

1.2. Cyclic Colourings of Embedded Graphs. — Given a surface S and a graph G embeddable in S , we denote by G^S that graph with a prescribed embedding in S . If the surface S is the sphere, we talk about a *plane graph* G^P . The *order* of a face of G^S is the number of vertices in its boundary; the maximum order of a face of G^S is denoted by $\Delta^*(G^S)$.

A *cyclic colouring* of an embedded graph G^S is a vertex colouring of G such that any two vertices in the boundary of the same face have distinct colours. The minimum number of colours required in a cyclic colouring of an embedded graph is called the *cyclic chromatic number* $\chi^*(G^S)$. This concept was introduced for plane graphs by Ore and Plummer [26], who also proved that for a plane graph G^P we have $\chi^*(G^P) \leq 2\Delta^*$. Borodin [3] (see also Jensen and Toft [14, page 37]) conjectured the following.

Conjecture 1.9 (Borodin [3]). — *For a plane graph G^P of maximum face order Δ^* we have $\chi^*(G^P) \leq \lfloor \frac{3}{2}\Delta^* \rfloor$.*

The bound in this conjecture is best possible. Consider the plane graph depicted in Figure 1(b): It has $3k$ vertices and has three faces of order $\Delta^* = 2k$. Since all pairs of vertices have a face they are both incident with, we need $3k = \lfloor \frac{3}{2}\Delta^* \rfloor$ colours in a cyclic colouring.

Borodin [3] also proved Conjecture 1.9 for $\Delta^* = 4$. For general values of Δ^* , the original bound $\chi^*(G^P) \leq 2\Delta^*$ of Ore and Plummer [26] was improved by Borodin *et al.* [6] to $\chi^*(G^P) \leq \lfloor \frac{9}{5} \Delta^* \rfloor$. The best known upper bound in the general case is due to Sanders and Zhao [29]: $\chi^*(G^P) \leq \lceil \frac{5}{3} \Delta^* \rceil$.

Although Wegner's and Borodin's Conjectures seem to be closely related, nobody has ever been able to bring to light a direct connection between them. Most of the results approaching these conjectures use the same ideas, but up until this point no one had proved a general theorem implying both a result on the colouring of the square and a result on the cyclic colouring of plane graphs (let alone on embedded graphs).

In order to show that our Theorem 1.2 provides an asymptotically best possible upper bound for the cyclic chromatic number for a graph G with some fixed embedding G^S , we need some extra notation. For each face f of G^S , add a vertex x_f . For any face f of G^S and any vertex v in the boundary of f , add an edge between v and x_f , and denote by G_F the graph obtained from G^S by this construction. Note that the vertex set of G_F consists of $V(G)$ and all the new vertices x_f , for f a face of G^S . Define a Σ -system Σ_F for G_F as follows: For each vertex $v \in V(G)$, let $\Sigma_F(v) = \emptyset$. For each vertex x_f , let $\Sigma_F(x_f)$ be all the neighbours of x_f . Observe that a (list) Σ_F -colouring of G_F colours the vertices of G in a way required for a cyclic (list) colouring of G^S , and that $\Delta(G_F; \Sigma_F) = \Delta^*(G^S)$.

(Note that in fact we have $\chi^*(G^S) \leq \chi(G_F, \Sigma_F) \leq \chi^*(G^S) + 1$. To get the second inequality, start with a cyclic colouring of G^S , add one extra colour, and colour all the vertices x_f with that colour. Similar inequalities hold for the list version.)

Using the upper bound on $\chi^*(G^S)$, we get the following corollary of Theorem 1.2.

Corollary 1.10. — *Let S be a fixed surface. Every embedding G^S of a graph G of maximum face order Δ^* has cyclic list chromatic number at most $\frac{3}{2} \Delta^* + o(\Delta^*)$.*

For an embedded graph G^S , the *cyclic clique number* $\omega^*(G^S)$ is the maximum size of a set $C \subseteq V$ such that every two vertices in C have some face they are both incident with. Note that the plane graph depicted in Figure 1(b) satisfies $\omega^*(G^P) = 3k = \lfloor \frac{3}{2} \Delta^* \rfloor$. This shows that the following corollary of Theorem 1.3 is best possible, up to the constant term.

Corollary 1.11. — *Let S be a fixed surface. Every embedded graph G^S of maximum face order Δ^* has cyclic clique number at most $\frac{3}{2} \Delta^* + O(1)$.*

For plane graphs, the proof of Theorem 1.3 guarantees that a plane graph G^P of maximum face order $\Delta^* \geq 11616$ has cyclic clique number at most $\frac{3}{2} \Delta^* + 76$.

2. Proof of Theorem 1.2

Our goal in this section is to show that for all surfaces S and all $\varepsilon > 0$, if we take β large enough (depending on S and ε), then for every graph $G = (V, E)$ embeddable in S , every choice of $\Sigma(v) \subseteq N_G(v)$ with $|\Sigma(v)| \leq \beta$ for all $v \in V$, and every assignment $L(v)$ of at least $(\frac{3}{2} + \varepsilon)\beta$ colours to all $v \in V$, there is a list Σ -colouring of G where each vertex receives a colour from its own list. In other words, we want an assignment $c(v)$ for each $v \in V$ such that:

- for all $v \in V$, we have $c(v) \in L(v)$;
- for all $u, v \in V$ with $uv \in E$, we have $c(u) \neq c(v)$; and
- for all $u, v \in V$ for which there is a $t \in V$ with $u, v \in \Sigma(t)$, we have $c(u) \neq c(v)$.

Before we present the actual proofs, we recall some of the important terminology, notation and facts concerning embeddings of graph in surfaces.

2.1. Graphs in Surfaces. — In this subsection, we give some background about graphs embedded in a surface. For more details, the reader is referred to [23]. Here, by a *surface* we mean a compact 2-dimensional surface without boundary. An *embedding* of a graph G in a surface S is a drawing of G on S so that all vertices are distinct, and every edge forms a simple arc connecting in S the vertices it joins, so that the interior of every edge is disjoint from other vertices and edges. A *face* of this embedding (or just a *face of G* , for short) is an arc-wise connected component of the space obtained by removing the vertices and edges of G from the surface S .

We say that an embedding is *cellular* if every face is homeomorphic to an open disc in \mathbb{R}^2 .

A surface can be orientable or non-orientable. The *orientable surface* \mathbb{S}_h of *genus h* is obtained by adding $h \geq 0$ ‘handles’ to the sphere; while the *non-orientable surface* \mathbb{N}_k of *genus k* is formed by adding $k \geq 1$ ‘cross-caps’ to the sphere. The *genus* $\mathbf{g}(G)$ and *non-orientable genus* $\tilde{\mathbf{g}}(G)$ of a graph G is the minimum h and the minimum k , resp., such that G has an embedding in \mathbb{S}_h , resp. in \mathbb{N}_k .

The following result will allow us to suppose that a graph G with known genus $\mathbf{g}(G)$ or non-orientable genus $\tilde{\mathbf{g}}(G)$ can be assumed to be embedded in a cellular way.

Lemma 2.1 ([23, Propositions 3.4.1 and 3.4.2]). — (i) *Every embedding of a connected graph G in $\mathbb{S}_{\mathbf{g}(G)}$ is cellular.*

(ii) *If G is a connected graph different from a tree, then there is an embedding of G in $\mathbb{N}_{\tilde{\mathbf{g}}(G)}$ that is cellular.*

The *Euler characteristic* $\chi(S)$ of a surface S is $2 - 2h$ if $S = \mathbb{S}_h$, and $2 - k$ if $S = \mathbb{N}_k$.

The basic result connecting all these concepts is *Euler’s Formula*: If G is a graph with an embedding in S , with vertex set V , edge set E and face set F , then

$$|V| - |E| + |F| \geq \chi(S).$$

Moreover, if the embedding is cellular, then we have equality in Euler’s Formula.

Finally, if v is a vertex of a graph G embedded in a surface S , then that embedding imposes two circular orders of the edges incident with v . Since we assume graphs to be simple, this corresponds to two circular orders of the neighbours of v . If S is orientable, then we can consistently choose one of the two clockwise orders for *all* vertices; if S is non-orientable, then such a choice is not possible. In our proofs that follow, it is not important that we can choose a consistent circular order; we only require that for each vertex v , there is at least one circular order of the neighbours around v .

If u_1, u_2 are consecutive neighbours of v (with respect to the chosen circular order), then there is a face that has the three vertices u_1, v, u_2 in its boundary. That immediately gives the following observation.

Lemma 2.2. — *Let G be a graph embedded in a surface S . Suppose u_1, u_2 are consecutive neighbours of v (with respect to the chosen circular order). Then the graph obtained by adding the edge u_1u_2 (if it is not already present) is still embeddable in S .*

That observation has the following corollary.

Lemma 2.3. — *Let G be a connected graph embedded in a surface S . If G has more than three vertices and is edge-maximal with respect to being embeddable in S , then every vertex has degree at least three.*

2.2. The First Steps. — For $P, Q \subseteq V$, the set of edges between P and Q is denoted by $E(P, Q)$, and the number of edges between P and Q is denoted by $e(P, Q)$ (edges with both ends in $P \cap Q$ are counted twice).

For a graph G with a Σ -system, and a vertex $v \in V$, we denote by $\sigma(v)$ the size of $\Sigma(v)$, i.e., $\sigma(v) = |\Sigma(v)|$. A Σ -neighbour of a vertex v is a vertex $u \neq v$ such that either u and v are adjacent, or there is some $t \in V$ with $u, v \in \Sigma(t)$. Denote the number of Σ -neighbours of v by $d^\Sigma(v)$. Note that we have

$$d^\Sigma(v) \leq d(v) + \sum_{t \text{ with } v \in \Sigma(t)} (\sigma(t) - 1).$$

An important tool in our proof of Theorem 1.2 is the following technical structural result, Lemma 2.4. Before stating this lemma, we need a few extra definitions. For an integer ζ , a *special ζ -pair* is a pair (X, Y) of disjoint subsets of vertices X and Y (possibly empty) with the following property:

- (i) *Every vertex in X has degree at least $\zeta + 1$. Every vertex $y \in Y$ has degree four, is adjacent to exactly two vertices of X , and the remaining neighbours of y have degree four as well.*

Given a special ζ -pair (X, Y) , for any $y \in Y$, let X^y be the set of two neighbours of y in X . For $W \subseteq X$, let Y^W be the set of all vertices $y \in Y$ with $X^y \subseteq W$ (that is, the set of vertices of Y having their two neighbours from X in W).

A special ζ -pair (X, Y) is called *very special* if in addition the following condition holds:

- (ii) *For all pairs of vertices $y, z \in Y$, if y and z are adjacent or have a common neighbour $w \notin X$, then $X^y = X^z$.*

The general structure of a very special ζ -pair is sketched in Figure 2.

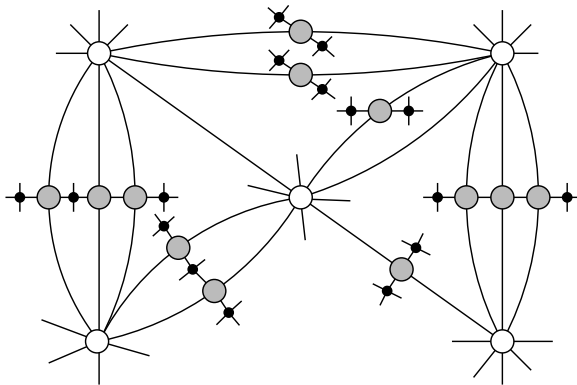


FIGURE 2. Sets X (white vertices) and Y (grey vertices) forming a very special 8-pair (X, Y) in a graph. Neighbours of vertices of Y not in X are depicted as small black vertices, and the remaining vertices are not depicted for the sake of clarity.

With these definitions, our structural lemma can be stated as follows:

Lemma 2.4. — *Let S be a fixed surface, set $\zeta_S^* = 132(3 - \chi(S))$, and let G be a graph embeddable in S . If G is edge-maximal with respect to being embeddable in S , then one of the following three properties holds.*

(S1) *Every vertex has degree at most ζ_S^* .*

(S2) *There is a vertex of degree at most five with at most one neighbour of degree more than ζ_S^* .*

(S3) *There exists a very special ζ_S^* -pair (X, Y) such that X, Y are both non-empty and for all non-empty subsets $W \subseteq X$, the following inequality holds:*

$$e(W, V \setminus Y) \leq e(W, Y \setminus Y^W) + \zeta_S^* |W|.$$

Very informally, Lemma 2.4 states that a graph that is maximally embeddable in some fixed surface, either contains one of two fairly simple configurations, or it contains a structure that internally satisfies a specific density-type condition.

Structure (S3) is at the heart of the above lemma. Although its description might appear technical at first sight, it will be clear later that it is the exact kind of density condition needed in the proofs of Theorems 1.2 and 1.3.

The proof of Lemma 2.4 can be found in Subsection 3.1. Observe that the value we use for ζ_S^* is probably far from best possible. The important point, to our mind, is that it only depends on (the Euler characteristic of) the surface S .

We continue with a description how to apply the lemma to prove Theorem 1.2. Suppose the theorem is false. Then there is a surface S and a real $\varepsilon > 0$ such that for every $\beta_{S,\varepsilon}$ we can find $\beta \geq \beta_{S,\varepsilon}$ and a graph G , together with a Σ -system of width at most β , such that $ch(G; \Sigma) > (\frac{3}{2} + \varepsilon)\beta$. Set $\zeta_S^* = 132(3 - \chi(S))$ and $\beta_S^* = \frac{2}{3}(\zeta_S^*)^2 = 11616(3 - \chi(S))^2$. Note that, as $\chi(S) \leq 2$, this means $\zeta_S^* \geq 132$ and $\beta_S^* \geq 11616$.

We start by assuming $\beta \geq \beta_S^*$; later (at the end of Subsection 2.3) we will add some further lower bounds for β that will depend on ε . With respect to this (yet to come) final choice of β , there exists a graph $G = (V, E)$ embeddable in S , together with a Σ -system of width at most β and a list-assignment L of at least $(\frac{3}{2} + \varepsilon)\beta$ colours to each vertex $v \in V$, such that G has no Σ -colouring from these lists. Choose such a graph G with the minimum number of vertices, and subject to this, with the maximum number of edges.

Certainly we can assume that G is connected (otherwise one of the components will be a smaller counterexample). Also, since each vertex has a list of more than $\frac{3}{2}\beta \geq 17424$ colours, G itself will have more than 17424 vertices.

Next we can assume that G is edge-maximal with respect to being embeddable in S . Otherwise we can add a new edge uv to G so that the resulting graph G_1 is still embeddable in S , and set $\Sigma_1 \equiv \Sigma$. It is clear that a list Σ_1 -colouring of G_1 is also a list Σ -colouring of G .

Fix some embedding of G in S . We continue with applying Lemma 2.4 to G .

2.2.1. The structure from (S1) is present in G . — This is the easiest case: If the degree of every vertex is at most ζ_S^* , then the number of Σ -neighbours of any vertex is at most $\zeta_S^* + \zeta_S^* \cdot (\zeta_S^* - 1) = (\zeta_S^*)^2$. But the number of colours in each list $L(v)$ is at least $(\frac{3}{2} + \varepsilon)\beta > \frac{3}{2}\beta_S^* = (\zeta_S^*)^2$. So a simple greedy colouring will do the job; contradicting that G is a counterexample.

2.2.2. The structure from (S2) is present in G . — So there is a vertex v of degree at most five, and at most one of its neighbours has degree more than ζ_S^* . By Lemma 2.3, since $|V| \geq 17424$, v has degree at least three. Hence it has a neighbour u of degree at most ζ_S^* . Form the

graph G_2 by contracting uv into a new vertex w (removing multiple edges if they appear). Set $V_2 = (V \setminus \{u, v\}) \cup \{w\}$. Let $\Sigma_2(w) = (\Sigma(u) \cup \Sigma(v)) \setminus \{u, v\}$. For a vertex $t \in V_2 \setminus \{w\}$, if $\Sigma(t)$ contains u , then set $\Sigma_2(t) = (\Sigma(t) \setminus \{u, v\}) \cup \{w\}$; otherwise set $\Sigma_2(t) = \Sigma(t) \setminus \{v\}$. Finally, give w the list of colours $L(w) = L(u)$. Note that G_2 is smaller than G and is still embeddable in S . Moreover, for every $t \in V_2 \setminus \{w\}$ we have $|\Sigma_2(t)| \leq |\Sigma(t)| \leq \beta$; while for w we have $|\Sigma_2(w)| \leq d_G(u) + d_G(v) \leq 5 + \zeta_S^* \leq \beta$.

So there exists a list Σ_2 -colouring of G_2 . We define a colouring of G as follows: Every vertex different from u and v keeps its colour from the colouring of G_2 . We give u the colour given to w in G_2 . Finally, we observe that for v we have

$$d^\Sigma(v) \leq d(v) + \sum_{t \text{ with } v \in \Sigma(t)} (\sigma(t) - 1) \leq 5 + 4(\zeta_S^* - 1) + (\beta - 1) = 4\zeta_S^* + \beta \leq \frac{3}{2}\beta,$$

since $\beta \geq \frac{2}{3}(\zeta_S^*)^2 \geq 8\zeta_S^*$. Since v has at least $(\frac{3}{2} + \varepsilon)\beta$ colours in its list, there exists a free colour for v , i.e., a colour different from the colour of all the vertices in conflict with v . We colour v with such a free colour. By the construction of G_2 and Σ_2 , it is easy to verify that this defines a list Σ -colouring of G , contradicting the choice of G as a counterexample.

2.2.3. The structure from (S3) is present in \mathbf{G} . — Let X and Y be two non-empty disjoint subsets of V such that the pair (X, Y) is a very special ζ_S^* -pair satisfying the condition of (S3). We can remove from X any vertex not adjacent to any vertex in Y .

Claim 2.5. — *For all $y \in Y$ we have that if $X^y = \{x_1, x_2\}$, then $y \in \Sigma(x_1) \cap \Sigma(x_2)$.*

Proof. — Suppose we have $y \notin \Sigma(x_1)$. Since (X, Y) is special, y has degree four, and it has a neighbour u not in X^y of degree four. We also have $d^\Sigma(y) \leq 4 + 2 \cdot (4 - 1) + (\sigma(x_2) - 1) \leq 9 + \beta$. By contracting the edge uy , we can argue similarly to Subsection 2.2.2 (with y now playing the role of v) to obtain a contradiction. \square

In the remainder of this subsection we describe how to reduce this case to a list edge-colouring problem. More precisely, we first define a modification of the original graph G into a smaller graph G_0 with vertex set $V \setminus Y$, inheriting a Σ -system from that of G , so that the minimality of G as a counterexample implies that G_0 admits a Σ -colouring. This colouring then provides a partial Σ -colouring of G , giving a colour to every vertex outside Y . In order to extend this partial colouring to the whole graph, we define a multigraph whose edges are indexed by the vertices in Y , so that an edge-colouring of that multigraph is exactly the extension of the Σ -colouring to Y we are looking for. In the next subsection we then describe how Kahn's approach to prove that the list chromatic index is asymptotically equal to the fractional chromatic index, can be used to conclude the proof of Theorem 1.2.

To define G_0 , we divide the vertices of Y into three parts according to their number of neighbours outside $X \cup Y$. Let Y' be the set of vertices in Y with no neighbour outside $X \cup Y$. Consider first the graph $G[V \setminus Y']$ induced on the set of vertices outside Y' . For each vertex $y \in Y'$, add an edge between its two neighbours $\{x_1, x_2\} = X^y$, if those are not already joined by an edge, and remove y from $\Sigma(x_1)$ and $\Sigma(x_2)$. Also, add x_1 to $\Sigma(x_2)$ and x_2 to $\Sigma(x_1)$. Note that after these changes, $\Sigma(x_1)$ and $\Sigma(x_2)$ cannot be larger than before (since, by Claim 2.5, $y \in \Sigma(x_i)$ for $i = 1, 2$).

For any vertex $y \in Y \setminus Y'$ with a unique neighbour u outside $X \cup Y$, contract the edge yu (removing multiple edges if they appear), and, by an abuse of the notation, call the new vertex u again. For the two vertices x_1 and x_2 in X^y , remove y from $\Sigma(x_1)$ and $\Sigma(x_2)$. For the vertex u itself, let $\Sigma(u)$ be equal to the set of all its neighbours.

And, finally, for any vertex $y \in Y \setminus Y'$ with exactly two neighbours u and u' outside $X \cup Y$, contract the edge yu (removing multiple edges if they appear), and, by an abuse of the notation, call the new vertex u again. For the two vertices x_1 and x_2 in X^y , remove y from $\Sigma(x_1)$ and $\Sigma(x_2)$. Add u to $\Sigma(u')$ and remove y from $\Sigma(u')$ (if it was in this set). For the vertex u itself, let $\Sigma(u)$ be equal to the set of all its neighbours. Note that u' has degree at most four in G , hence certainly $|\Sigma(u')| \leq \beta$.

The graph obtained after the modifications described above is denoted by G_0 , and the resulting sets by $\Sigma_0(v)$, $v \in V(G_0)$. Note that, by our abuse of the notation, G_0 has the vertex set $V_0 = V \setminus Y$. Next we observe that a vertex u of G outside $X \cup Y$ that was adjacent to a vertex $y \in Y$ (and hence may have been involved in one or more contractions) has degree four in G . Since vertices in Y have degree four as well, each contraction increases the degree by at most two. So in G_0 , such a vertex u has degree at most twelve, hence we certainly have $|\Sigma_0(u)| \leq \beta$. By the construction above, we saw that for every other vertex $v \in V_0$, we also have $|\Sigma_0(v)| \leq |\Sigma(v)| \leq \beta$ or $|\Sigma_0(v)| \leq d_{G_0}(v) \leq \beta$.

By its construction, G_0 is embeddable in S . Also by construction, and the remarks above, it is easy to verify the following statement.

Claim 2.6. — *If $u, v \in V_0$ are adjacent in G , then u, v are adjacent in G_0 . If $u, v \in V_0$ and there is a $t \in V$ with $u, v \in \Sigma(t)$, then u, v are either adjacent in G_0 , or there is a $t_0 \in V_0$ with $u, v \in \Sigma(t_0)$.*

For each vertex $v \in V_0$ set $L_0(v) = L(v)$. Since $Y \neq \emptyset$, by the minimality of G , the graph G_0 admits a list Σ_0 -colouring c_0 with respect to the list assignment L_0 .

We now transform this colouring into a partial list Σ -colouring of G with respect to the original list assignment L , by just setting $c(v) = c_0(v)$ for each vertex $v \in V_0 = V \setminus Y$. By Claim 2.6, this is indeed a good partial Σ -colouring of all the vertices of $V \setminus Y$ in G . The difficult part of the proof is to show that c can be extended to Y .

By assumption, at the beginning every vertex in Y has a list of at least $(\frac{3}{2} + \varepsilon)\beta$ available colours. For each vertex y in Y , let us remove from $L(y)$ the colours which are forbidden for y according to the partial Σ -colouring c of G . In the worst case, these forbidden colours are exactly the colours of the vertices of $V \setminus Y$ at distance at most two from y .

Let us define the multigraph H as follows: H has vertex set X . And for each vertex $y \in Y$ we add an edge e_y between the two neighbours of y in X (in other words, between the two vertices in X^y). Note that this process may produce multiple edges. We associate a list $L(e_y)$ to e_y in H by taking the list of y obtained after removing the set of forbidden colours for y from the original list $L(y)$.

In what follows, following the usual terminology for multigraphs, we denote by $d_H(x)$ the degree of the vertex x in the multigraph H , i.e., the number of edges incident with x in H . By Claim 2.5 we have $N_G(x) \cap Y \subseteq \Sigma(x)$ for every $x \in X$, which guarantees $d_H(x) \leq \sigma(x)$.

We now prove the following lemma.

Lemma 2.7. — *A list edge-colouring for H , with the list assignment L defined as above, provides an extension of c to a list Σ -colouring of G by giving to each vertex $y \in Y$ the colour of the edge e_y in H .*

Proof. — This follows since the pair (X, Y) is very special: For every two vertices $y, z \in Y$, if y and z are adjacent or have a common neighbour $w \notin X$, then $X^y = X^z$. This proves that the two vertices adjacent in Y or with a common neighbour not in X define parallel

edges in H and so will have different colours. If two vertices y_1 and y_2 of Y have a common neighbour in X , e_{y_1} and e_{y_2} will be adjacent in H and so will get different colours. Since we have already removed from the list of vertices in Y the set of forbidden colours (defined by the colours of the vertices in $V \setminus Y$), there will be no conflict between the colours of a vertex in Y and a vertex in $V \setminus Y$. We conclude that the edge-colouring of H will provide an extension of c to a list Σ -colouring of G . \square

The following lemma provides a lower bound on the size of $L(e)$ for the edges e in H .

Lemma 2.8. — *Let $e = x_1x_2$ be an edge in H . Then we have*

$$|L(e)| \geq \left(\frac{3}{2} + \varepsilon\right) \beta - (\sigma(x_1) - d_H(x_1)) - (\sigma(x_2) - d_H(x_2)) - 10.$$

Proof. — Let y be the vertex in Y such that $e = e_y$. By the definition of H , $X^y = \{x_1, x_2\}$. Let Z be the set of vertices in $V \setminus X$ adjacent to y in G . Then, since (X, Y) is a special ζ_S^* -pair, $|Z| \leq 2$ and $|N_G(Z) \setminus Y| \leq 6$. The colours that are possibly forbidden for y are the colours of $\{x_1, x_2\}$, plus the colours of vertices in $(Z \cup N_G(Z)) \setminus Y$, plus the colours of vertices in $(\Sigma(x_1) \setminus Y) \cup (\Sigma(x_2) \setminus Y)$ (note that these colours all come from the vertices outside Y). The number of vertices in these three sets add up to at most $10 + (\sigma(x_1) - d_H(x_1)) + (\sigma(x_2) - d_H(x_2))$. The lemma follows. \square

We finish this subsection by applying Lemma 2.4 in order to obtain information on the density of subgraphs in H , which we will need in the next subsection. Recall that for all non-empty subsets $W \subseteq X$, Y^W denotes the set of vertices $y \in Y$ with $X^y \subseteq W$ (that is, the set of vertices of Y having their two neighbours from X in W). By (S3) we have for all non-empty $W \subseteq X$,

$$e_G(W, V \setminus Y) \leq e_G(W, Y \setminus Y^W) + \zeta_S^* |W|.$$

This inequality has the following interpretation in H .

Lemma 2.9. — *For all non-empty subsets $W \subseteq X (= V(H))$, we have*

$$\sum_{w \in W} (\sigma(w) - d_H(w)) \leq e_H(W, X \setminus W) + \zeta_S^* |W|.$$

Proof. — First note that $\sum_{w \in W} (\sigma(w) - d_H(w)) \leq \sum_{w \in W} (d_G(w) - d_H(w)) = e_G(W, V \setminus Y)$.

We also have $e_G(W, Y \setminus Y^W) = e_H(W, X \setminus W)$. Combining these two observations with the formula in (S3) immediately gives the required inequality. \square

At this point, our aim will be to apply Kahn's approach to the multigraph H with the list assignment L , to prove the existence of a proper list edge-colouring for H . This is described in the next subsection.

We summarise the properties we assume are satisfied by the multigraph H and the list assignment L of the edges of H . For these conditions we just consider $\sigma(v)$ as an integer with certain properties, assigned to each vertex v of H .

(H1) For all vertices v in H , we have $d_H(v) \leq \sigma(v) \leq \beta$.

(H2) For all edges $e = uv$ in H , $|L(e)| \geq \left(\frac{3}{2} + \varepsilon\right) \beta - (\sigma(u) - d_H(u)) - (\sigma(v) - d_H(v)) - 10$.

(H3) For all non-empty subsets $W \subseteq V(H)$, $\sum_{w \in W} (\sigma(w) - d_H(w)) \leq e_H(W, V(H) \setminus W) + \zeta_S^* |W|$,

for some constant ζ_S^* .

2.3. The Matching Polytope and Edge-Colourings. — We briefly describe the matching polytope of a multigraph. More about this subject can be found in [30, Chapter 25].

Let H be a multigraph with m edges. Let $\mathcal{M}(H)$ be the set of all matchings of H , including the empty matching. For each $M \in \mathcal{M}(H)$, let us define the m -dimensional characteristic vector $\mathbf{1}_M$ as follows: $\mathbf{1}_M = (x_e)_{e \in E(H)}$, where $x_e = 1$ for an edge $e \in M$, and $x_e = 0$ otherwise. The matching polytope of H , denoted $\mathcal{MP}(H)$, is the polytope defined by taking the convex hull of all the vectors $\mathbf{1}_M$ for $M \in \mathcal{M}(H)$. Also, for any real number λ , we set $\lambda \mathcal{MP}(H) = \{ \lambda x \mid x \in \mathcal{MP}(H) \}$.

Edmonds [9] gave the following characterisation of the matching polytope.

Theorem 2.10 (Edmonds [9]). — *A vector $\vec{x} = (x_e)$ is in $\mathcal{MP}(H)$ if and only if $x_e \geq 0$ for all x_e and the following two types of inequalities are satisfied:*

- For all vertices $v \in V(H)$, $\sum_{e: v \text{ incident to } e} x_e \leq 1$;
- for all subsets $W \subseteq V(H)$ with $|W| \geq 3$ and $|W|$ odd, $\sum_{e \in E(W)} x_e \leq \frac{1}{2}(|W| - 1)$.

The significance of the matching polytope and its relation to list edge-colouring is indicated by the following important result.

Theorem 2.11 (Kahn [17]). — *For all real numbers δ, ν , $0 < \delta < 1$ and $\nu > 0$, there exists a $\Delta_{\delta, \nu}$ such that for all $\Delta \geq \Delta_{\delta, \nu}$ the following holds. If H is a multigraph and L is a list assignment of colours to the edges of H so that*

- H has maximum degree at most Δ ;
- for all edges $e \in E(H)$, $|L(e)| \geq \nu \Delta$;
- the vector $\vec{x} = (x_e)$ with $x_e = \frac{1}{|L(e)|}$ for all $e \in E(H)$ is an element of $(1 - \delta) \mathcal{MP}(H)$.

Then there exists a proper edge-colouring of H where each edge gets a colour from its own list.

The theorem above is actually not explicitly stated this way in [17], but can be obtained from the appropriate parts of that paper. We give some further details about this in the final section of this paper.

The next lemma allows us to use Theorem 2.11 to complete the proof.

Lemma 2.12. — *Let β and ζ be positive real numbers. Let H be a multigraph with a map $\sigma : V(H) \rightarrow \mathbb{N}$, and a weighting $(b_e)_{e \in E(H)}$ of the edges with positive real numbers satisfying the following three conditions:*

- (H1') For all vertices v in H , $d_H(v) \leq \sigma(v) \leq \beta$.
- (H2') For all edges $e = uv$ in H , $b_e \geq \left(\frac{3}{2}\beta + \frac{9}{2}\zeta\right) - (\sigma(u) - d_H(u)) - (\sigma(v) - d_H(v))$.
- (H3') For all non-empty $W \subseteq V(H)$, $\sum_{w \in W} (\sigma(w) - d_H(w)) \leq e_H(W, V(H) \setminus W) + \zeta |W|$.

Then for all edges $e \in E(H)$, we have $b_e \geq \frac{1}{2}\beta$. And the vector $(1/b_e)_{e \in E(H)}$ is in $\mathcal{MP}(H)$.

The proof of Lemma 2.12 will be given in Subsection 3.2. This lemma guarantees that for all $\varepsilon > 0$, there exists a β_ε such that for all $\beta \geq \beta_\varepsilon$, Theorem 2.11 can be applied to a multigraph H with an edge list assignment L satisfying properties (H1)–(H3) stated at the end of the previous subsection.

To see this, take $\delta_\varepsilon = \frac{\varepsilon}{3+2\varepsilon}$, so $0 < \delta_\varepsilon < 1$. In order to be able to apply Theorem 2.11, we want to prove the existence of $\beta_{\varepsilon, \zeta_S^*}$ such that for any $\beta \geq \beta_{\varepsilon, \zeta_S^*}$, the vector $\vec{x} = (x_e)$, $x_e = \frac{1}{|L(e)|}$, is in $(1 - \delta_\varepsilon) \mathcal{MP}(H)$. Let ζ_S^* be the constant described in condition (H3). By condition (H2), we have for all $e = uv$ in H ,

$$\begin{aligned} (1 - \delta_\varepsilon) |L(e)| &\geq (1 - \delta_\varepsilon) \left(\left(\frac{3}{2} + \varepsilon \right) \beta - (\sigma(u) - d_H(u)) - (\sigma(v) - d_H(v)) - 10 \right) \\ &\geq (1 - \delta_\varepsilon) \left(\frac{3}{2} + \varepsilon \right) \beta - (\sigma(u) - d_H(u)) - (\sigma(v) - d_H(v)) - 10 \\ &= \left(\frac{3}{2} \beta + \frac{1}{2} \varepsilon \beta \right) - (\sigma(u) - d_H(u)) - (\sigma(v) - d_H(v)) - 10. \end{aligned}$$

Let $\beta_{\varepsilon, \zeta_S^*} = \frac{9\zeta_S^* + 20}{\varepsilon}$. For $\beta \geq \beta_{\varepsilon, \zeta_S^*}$ we have

$$(1 - \delta_\varepsilon) |L(e)| \geq \left(\frac{3}{2} \beta + \frac{9}{2} \zeta_S^* \right) - (\sigma(u) - d_H(u)) - (\sigma(v) - d_H(v)).$$

So by Lemma 2.12, taking $b_e = (1 - \delta_\varepsilon) |L(e)|$, the vector $\left(\frac{x_e}{1 - \delta_\varepsilon} \right)_{e \in E(H)}$ is in $\mathcal{MP}(H)$. We infer that $\vec{x} \in (1 - \delta_\varepsilon) \mathcal{MP}(H)$.

Now set $\beta_{S, \varepsilon} = \max\{\beta_S^*, \beta_{\varepsilon, \zeta_S^*}, \Delta_{\delta_\varepsilon, 1/2}\}$ (where β_S^*, ζ_S^* are determined by Lemma 2.4, δ_ε and $\beta_{\varepsilon, \zeta_S^*}$ are given above, and $\Delta_{\delta_\varepsilon, 1/2}$ is according to Theorem 2.11), and assume $\beta \geq \beta_{S, \varepsilon}$. Then, using Lemma 2.12, we can apply Theorem 2.11 which implies that the multigraph H defined in Subsection 2.2 has a list edge-colouring corresponding to the list assignment L . Lemma 2.7 then implies that the colouring c can be extended to a list Σ -colouring of the original graph G . This final contradiction completes the proof of Theorem 1.2. \blacksquare

3. Proofs of the Main Lemmas

We use the terminology and notation from the previous sections.

3.1. Proof of Lemma 2.4. — Let S be a surface, set $\zeta_S^* = 132(3 - \chi(S))$, and let G be a graph embeddable in S , so that G is edge-maximal with respect to being embeddable in S .

From Lemma 2.2, we immediately obtain the following.

Claim 3.1. — *For any vertex v and any two consecutive neighbours u_1, u_2 of v (consecutive with respect to the chosen circular order imposed by the embedding), we have $u_1 u_2 \in E(G)$.*

Next we prove that we can assume G has a cellular embedding in S . If G is a tree, then every leaf will give a structure from (S2). So we can assume G is not a tree. Assume S is orientable with genus h . By the definition of $\mathbf{g}(G)$, we must have $\mathbf{g}(G) \leq h$, and hence $\chi(\mathbb{S}_{\mathbf{g}(G)}) = 2 - 2\mathbf{g}(G) \geq 2 - 2h = \chi(S)$. That also means that the constant in Lemma 2.4 satisfies $\zeta_{\mathbb{S}_{\mathbf{g}(G)}}^* \leq \zeta_S^*$. Hence if we prove the lemma assuming G is embeddable in $\mathbb{S}_{\mathbf{g}(G)}$, then the lemma for G embeddable in S directly follows. So we can use Lemma 2.4 with the surface $\mathbb{S}_{\mathbf{g}(G)}$ instead of S , and by Lemma 2.1, we can use a cellular embedding of G in $\mathbb{S}_{\mathbf{g}(G)}$.

If S is non-orientable, then exactly the same argument can be applied, this time using the surface $\mathbb{N}_{\mathbf{g}(G)}$ (and using the assumption that G is not a tree).

We need some further notation and terminology. The set of faces of G is denoted by F . Recall that since the embedding in S is cellular, every face is homeomorphic to an open disk in \mathbb{R}^2 . For such a face f , a *boundary walk of f* is a walk consisting of vertices and edges as they are encountered when walking along the whole boundary of f , starting at some vertex.

The *degree of a face* f , denoted $d(f)$, is the number of edges on the boundary walk of f . Note that this means that some edges may be counted more than once. The *order* of a face is the number of vertices in its boundary. We always have that the order of f is at most $d(f)$.

Now suppose that G does not contain any of the structures (S1) or (S2). In order to prove Lemma 2.4, we only need to prove that G contains structure (S3). In other words, we need to prove that G contains a very special ζ_S^* -pair (X, Y) with X and Y non-empty which satisfies the inequality of (S3) for all non-empty subsets $W \subseteq X$.

We easily see that G has at least $\zeta_S^* + 2 \geq 134$ vertices (otherwise it contains structure (S1)). So by Lemma 2.3 we know that all vertices have degree at least three.

Let us call the vertices of degree at least $\zeta_S^* + 1$ *big*; the other vertices are called *small*. We use B to denote the set of big vertices.

Since we assumed that G does not contain structure (S2), we immediately get:

Claim 3.2. — *All vertices of degree at most five have at least two big neighbours.*

We continue our analysis using the classical technique of *discharging*. Give each vertex v an initial charge $\rho(v) = 6d(v) - 36$. Since G is simple and has a cellular embedding in S , every face has degree at least three. This gives $2|E| \geq 3|F|$, and hence, by Euler's Formula, $\sum_{v \in V} \rho(v) = 12|E| - 36|V| \leq -36|V| + 36|E| - 36|F| = -36\chi(S)$.

We further redistribute charges according to the following rules:

- (R1) Each vertex of degree three that is adjacent to three big vertices receives a charge 6 from each of its neighbours.
- (R2) Each vertex of degree three that is adjacent to two big vertices receives a charge 9 from each of its big neighbours.
- (R3) Each vertex of degree four that is adjacent to four big vertices receives a charge 3 from each of its big neighbours.
- (R4) Each vertex of degree four that is adjacent to three big vertices receives a charge 4 from each of its big neighbours.
- (R5) Each vertex of degree four that is adjacent to two big vertices receives a charge 6 from each of its big neighbours.
- (R6) Each vertex of degree five receives a charge 3 from each of its big neighbours.

Denote the resulting charge of a vertex $v \in V$ after applying rules (R1) – (R6) by $\rho'(v)$. Since the global charge has been preserved, we have $\sum_{v \in V} \rho'(v) \leq -36\chi(S)$. We will show that for most $v \in V$, $\rho'(v)$ is non-negative.

Combining Claim 3.2 with rules (R1) – (R6) and our knowledge that $\rho(v) = 6d(v) - 36$, we find that $\rho'(v) = 0$ if $d(v) = 3, 4$, while $\rho'(v) \geq 0$ if $d(v) = 5$. If v is a small vertex with $d(v) \geq 6$, we have $\rho'(v) = \rho(v) = 6d(v) - 36 \geq 0$.

It follows that we must have

$$(1) \quad \sum_{v \in B} \rho'(v) \leq -36\chi(S).$$

To derive the relevant consequence of that formula, we must make a detailed analysis of the neighbours of vertices in B .

As we explained in Subsection 2.1, the embedding of G in S allows us to choose a circular order on the neighbours of each vertex v . By Claim 3.1 we know that two consecutive vertices

in this order are adjacent. If u is a neighbour of v , then by u^+, u^{++} we denote the successor and second successor of u in the circular order of neighbours of v , while u^-, u^{--} denote the predecessor and second predecessor of u in that order.

We distinguish five different types of neighbours of a vertex $v \in B$:

$$\begin{aligned} M_1(v) &= \{u \in N(v) \mid \{u^-, u^{--}, u^+, u^{++}\} \cap B \neq \emptyset\}; \\ M_{4a}(v) &= \{u \in N(v) \setminus M_1(v) \mid d(u) = 4 \text{ and } u^- \text{ or } u^+ \text{ have degree at least five}\}; \\ M_{4b}(v) &= \{u \in N(v) \setminus M_1(v) \mid d(u) = d(u^-) = d(u^+) = 4\}; \\ M_5(v) &= \{u \in N(v) \setminus M_1(v) \mid d(u) = 5\}; \\ M_6(v) &= \{u \in N(v) \setminus M_1(v) \mid d(u) \geq 6\}. \end{aligned}$$

First observe that if a neighbour u of v has degree three, then u^- or u^+ is in B . This follows since by Claim 3.1, u^- and u^+ are also neighbours of u . And by Claim 3.2, a vertex of degree three has at least two big neighbours. From this observation we also get that if $u \in N(v) \setminus M_1(v)$ is a small vertex, then u^- and u^+ both have degree at least four.

As a consequence, every neighbour of v is in exactly one set. Our aim in the following, in order to prove Lemma 2.4, is to show that most neighbours of vertices $v \in B$ are in $M_{4b}(v)$.

We now evaluate the charge that a vertex $v \in B$ has given to its neighbours. If $u \in M_1(v)$, then v gave at most $9 + 9 + 9 = 27$ to $\{u^-, u, u^+\}$; if $u \in M_{4a}(v)$, then v gave at most $3 + 6 + 6 = 15$ to $\{u^-, u, u^+\}$; if $u \in M_{4b}(v)$, then v gave at most $6 + 6 + 6 = 18$ to $\{u^-, u, u^+\}$; if $u \in M_5(v)$, then v gave at most $6 + 3 + 6 = 15$ to $\{u^-, u, u^+\}$; and, finally, if $u \in M_6(v)$, then v gave at most $6 + 0 + 6 = 12$ to $\{u^-, u, u^+\}$. Setting $m_1 = |M_1(v)|$, $m_{4a} = |M_{4a}(v)|$, $m_{4b} = |M_{4b}(v)|$, $m_5 = |M_5(v)|$, and $m_6 = |M_6(v)|$, we can conclude that v gave at most

$$\begin{aligned} &\frac{1}{3}(27m_1 + 15m_{4a} + 18m_{4b} + 15m_5 + 12m_6) \\ &\leq 9m_1 + 6m_{4b} + 5(m_{4a} + m_5 + m_6) \leq 5d(v) + 4m_1 + m_{4b} \end{aligned}$$

to its neighbourhood. This means that the remaining charge $\rho'(v)$ of a vertex $v \in B$ must satisfy

$$\rho'(v) \geq (6d(v) - 36) - (5d(v) + 4m_1 + m_{4b}) = d(v) - m_{4b} - 4m_1 - 36.$$

By definition, $|M_1(v)|$ is at most four times the number of neighbours of v in B . Consider the subgraph $G[B]$ of G induced by B . As a subgraph of G , this graph is embeddable in S . As it is simple as well, no face of such an embedding is incident with two or fewer edges. So Euler's Formula means that $G[B]$ has at most $3|B| - 3\chi(S)$ edges, and hence

$$\sum_{v \in B} |M_1(v)| \leq \sum_{v \in B} 4d_{G[B]}(v) = 8|E(G[B])| \leq 24|B| - 24\chi(S).$$

Combining the last two inequalities with (1) gives

$$-36\chi(S) \geq \sum_{v \in B} \rho'(v) \geq \sum_{v \in B} (d(v) - |M_{4b}(v)|) - 4(24|B| - 24\chi(S)) - 36|B|.$$

Using that $B \neq \emptyset$ (otherwise G contains structure (S1)) and $\chi(S) \leq 2$, this can be written as

$$\sum_{v \in B} (d(v) - |M_{4b}(v)|) \leq 132|B| - 132\chi(S) < 132|B| + 132(2 - \chi(S)) \leq 132(3 - \chi(S))|B|.$$

Define $X_0 = B$ and $Y_0 = \bigcup_{v \in B} M_{4b}(v)$. Note that the previous inequality can be written

$$(2) \quad e(X_0, V \setminus Y_0) < \zeta_S^* |X_0|.$$

Also observe that the pair (X_0, Y_0) is a special ζ_S^* -pair: The vertices in X_0 are the big vertices, hence have degree at least $\zeta_S^* + 1$. For all vertices $u \in Y$ we have $u \in M_{4b}(v)$ for some $v \in B$, and hence u, u^- and u^+ have degree four in G , and the fourth neighbour of u is in $B = X_0$ by Claim 3.2.

We need some more information about the neighbours of vertices in Y_0 .

Claim 3.3. — *Let v be a big vertex, $u \in M_{4b}(v)$, and w be the big neighbour of u different from v . Then all of vu^+, vu^-, wu^+ and wu^- are edges of G .*

Proof. — Consider the circular order of the neighbours of u imposed by the embedding. In any circular order different from (v, u^+, w, u^-) or the reverse, u^+ and u^- are consecutive. By Claim 3.1, this means that $u^+u^- \in E$. So the neighbours of u^- are $\{v, u, u^+, u^{--}\}$. Since $u, u^+, u^{--} \notin B$, by the definition of $M_{4b}(v)$, that means u^- has only one big neighbour, contradicting Claim 3.2.

So the only possible circular orders are (v, u^+, w, u^-) or the reverse, and the result follows by Claim 3.1. \square

Using Claim 3.3, it follows easily that if $y, z \in Y_0$ are adjacent, then $X_0^y = X_0^z$; while if $y, z \in Y_0$ share a neighbour $u \notin X_0$, then u has degree four and its two neighbours distinct from y and z are in X_0^y and in X_0^z . This gives $X_0^y = X_0^z$.

Thus, we have shown that the pair (X_0, Y_0) is very special.

Since X_0 and Y_0 are non-empty, we are done if the pair (X_0, Y_0) also satisfies the inequalities of (S3) for any non-empty subset $W \subseteq X_0$. Suppose this is not the case. So there must exist a set $Z_1 \subseteq X_0$ with

$$e(Z_1, V \setminus Y_0) > e(Z_1, Y_0 \setminus Y_0^{Z_1}) + \zeta_S^* |Z_1|.$$

Define $X_1 = X_0 \setminus Z_1$ and $Y_1 = Y_0^{X_1}$. Again, by construction, it is easy to see that (X_1, Y_1) is a very special ζ_S^* -pair. If it does not satisfy condition (S3), we iterate the process (see Figure 3) and eventually obtain a very special ζ_S^* -pair (X_k, Y_k) satisfying condition (S3). To conclude the proof, we only need to check that X_k and Y_k are non-empty.

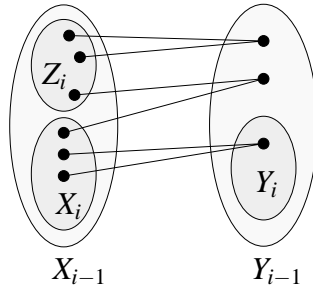


FIGURE 3. $X_i = X_{i-1} \setminus Z_i$ and $Y_i = Y_{i-1}^{X_i}$.

Let $1 \leq i \leq k$. Since $X_i = X_{i-1} \setminus Z_i$, we have

$$\begin{aligned} e(X_i, V \setminus Y_i) &= e(X_{i-1}, V \setminus Y_i) - e(Z_i, V \setminus Y_i) \\ &= e(X_{i-1}, V \setminus Y_{i-1}) + e(X_{i-1}, Y_{i-1} \setminus Y_i) - e(Z_i, V \setminus Y_{i-1}) - e(Z_i, Y_{i-1} \setminus Y_i) \\ &= e(X_{i-1}, V \setminus Y_{i-1}) - e(Z_i, V \setminus Y_{i-1}) + e(X_i, Y_{i-1} \setminus Y_i). \end{aligned}$$

Since $Y_i = Y_{i-1}^{X_i}$, every neighbour $u \in Y_{i-1} \setminus Y_i$ of a vertex in X_i has exactly one neighbour in Z_i (see Figure 3). Hence, $e(X_i, Y_{i-1} \setminus Y_i) = e(Z_i, Y_{i-1} \setminus Y_{i-1}^{Z_i})$. So we have

$$e(X_{i-1}, V \setminus Y_{i-1}) = e(X_i, V \setminus Y_i) + e(Z_i, V \setminus Y_{i-1}) - e(Z_i, Y_{i-1} \setminus Y_{i-1}^{Z_i}).$$

By the definition of Z_i , we have $e(Z_i, V \setminus Y_{i-1}) > e(Z_i, Y_{i-1} \setminus Y_{i-1}^{Z_i}) + \zeta_S^* |Z_i|$. Combining the last two expressions gives

$$e(X_{i-1}, V \setminus Y_{i-1}) > e(X_i, V \setminus Y_i) + \zeta_S^* |Z_i|.$$

Setting $Z^* = \bigcup_{1 \leq i \leq k} Z_i$, we have $e(X_k, V \setminus Y_k) < e(X_0, V \setminus Y_0) - \zeta_S^* |Z^*|$. As a consequence, using (2),

$$|Z^*| < \frac{e(X_0, V \setminus Y_0) - e(X_k, V \setminus Y_k)}{\zeta_S^*} \leq \frac{e(X_0, V \setminus Y_0)}{\zeta_S^*} < \frac{\zeta_S^* |X_0|}{\zeta_S^*} = |X_0|.$$

Since $X_k = X_0 \setminus Z^*$, this implies $|X_k| > 0$, which leads to $X_k \neq \emptyset$.

Finally, let $v \in X_k \neq \emptyset$ and assume $Y_k = \emptyset$. Taking $W = \{v\}$ in the inequality in (S3) (which by construction is satisfied by (X_k, Y_k)), we obtain $d(v) \leq \zeta_S^*$. Since v is a big vertex, $d(v) \geq \zeta_S^* + 1$. This contradiction means that we must have $Y_k \neq \emptyset$, which concludes the proof of Lemma 2.4. \blacksquare

3.2. Proof of Lemma 2.12. — We recall the hypotheses of the lemma: We have positive real numbers β and ζ ; H is a multigraph; each vertex v of H has an associated integer $\sigma(v)$; and for each edge e a positive real number b_e is given. In this subsection, all degrees $d(v)$ are in the multigraph H .

The following three conditions are satisfied:

(H1') For all vertices v in H , $d(v) \leq \sigma(v) \leq \beta$.

(H2') For all edges $e = uv$ in H , $b_e \geq \left(\frac{3}{2}\beta + \frac{9}{2}\zeta\right) - (\sigma(u) - d(u)) - (\sigma(v) - d(v))$.

(H3') For all non-empty subsets $W \subseteq V(H)$, $\sum_{w \in W} (\sigma(w) - d(w)) \leq e_H(W, V(H) \setminus W) + \zeta |W|$.

In the proof that follows, we will show that the vector $\vec{x} = (x_e)$, $x_e = 1/b_e$, is in $\mathcal{MP}(H)$.

For an edge $e = uv$ in H , define

$$(3) \quad a_e = \left(\frac{3}{2}\beta + \frac{9}{2}\zeta\right) - (\sigma(u) - d(u)) - (\sigma(v) - d(v)) \quad \text{and} \quad y_e = \frac{1}{a_e}.$$

We will in fact prove that the vector $\vec{y} = (y_e)$ is in the matching polytope $\mathcal{MP}(H)$. Since $b_e \geq a_e$, we have $x_e = 1/b_e \leq 1/a_e = y_e$. So, by Edmonds' characterisation of the matching polytope, if $\vec{y} \in \mathcal{MP}(H)$, this guarantees that $\vec{x} \in \mathcal{MP}(H)$, as required.

Applying condition (H3') to the set $W = \{v\}$ gives $\sigma(v) - d(v) \leq d(v) + \zeta$, which implies:

(a) For all vertices $v \in V(H)$, we have $d(v) \geq \frac{1}{2}(\sigma(v) - \zeta)$.

Let $e = uv$ be an edge of H . If we use the estimate above for both u and v in the definition of a_e in (3), and recalling that $\sigma(u), \sigma(v) \leq \beta$, we obtain

$$a_e \geq \frac{3}{2}\beta + \frac{9}{2}\zeta - \frac{1}{2}\sigma(u) - \frac{1}{2}\sigma(v) - \zeta \geq \frac{1}{2}\beta + \frac{7}{2}\zeta.$$

On the other hand, if we use observation (a) for u only, we get

$$a_e \geq d(v) + \frac{3}{2}\beta + \frac{9}{2}\zeta - \frac{1}{2}\sigma(u) - \sigma(v) - \frac{1}{2}\zeta \geq d(v) + 4\zeta.$$

Hence, the following two conclusions hold.

(b) For all edges $e = uv$ in $E(H)$, we have $a_e \geq d(v) + 4\zeta$.

(c) For all edges $e \in E(H)$, we have $a_e \geq \frac{1}{2}\beta + \frac{7}{2}\zeta$.

Note that observation (c) also gives $b_e \geq a_e \geq \frac{1}{2}\beta$ for all $e \in E(H)$, as required.

By observation (b), we find, since $\zeta > 0$,

$$\sum_{e \ni v} \frac{1}{a_e} \leq d(v) \cdot \frac{1}{d(v) + 4\zeta} < 1,$$

which shows that

Claim 3.4. — For all vertices $v \in V(H)$, we have $\sum_{e \ni v} y_e < 1$.

Using Theorem 2.10, all that remains is to prove that for all $W \subseteq V(H)$ with $|W| \geq 3$ and $|W|$ odd, we have $\sum_{e \in E(W)} y_e \leq \frac{1}{2}(|W| - 1)$. We will actually prove this for all $|W| \geq 3$. Note that

we can certainly assume $E(W) \neq \emptyset$.

Using observation (b), we infer that

$$\sum_{e \in E(W)} \frac{1}{a_e} \leq \frac{1}{2} \sum_{u \in W} \frac{d_{H[W]}(u)}{d(u) + 4\zeta} = \frac{1}{2} \sum_{u \in W} \left(\frac{d(u)}{d(u) + 4\zeta} - \frac{d(u) - d_{H[W]}(u)}{d(u) + 4\zeta} \right).$$

Since $\frac{d(u)}{d(u) + 4\zeta} \leq \frac{\beta}{\beta + 4\zeta}$ and $\frac{d(u) - d_{H[W]}(u)}{d(u) + 4\zeta} \geq \frac{d(u) - d_{H[W]}(u)}{\beta + 4\zeta}$, this implies

$$\sum_{e \in E(W)} \frac{1}{a_e} \leq \frac{1}{2} |W| \frac{\beta}{\beta + 4\zeta} - \frac{1}{2} \frac{e(W, W^c)}{\beta + 4\zeta}.$$

Here we used that $\sum_{u \in W} (d(u) - d_{H[W]}(u)) = e(W, W^c)$, where $W^c = V(H) \setminus W$.

If $e(W, W^c) \geq \beta$, we obtain, since $\zeta > 0$,

$$\sum_{e \in E(W)} y_e \leq \frac{1}{2}(|W| - 1) \cdot \frac{\beta}{\beta + 4\zeta} < \frac{1}{2}(|W| - 1).$$

So we can assume in the following that $e(W, W^c) \leq \beta$, in which case Condition (H3') of Lemma 2.12 implies

$$\sum_{u \in W} (\sigma(u) - d(u)) \leq e(W, W^c) + \zeta |W| \leq \beta + \zeta |W|.$$

For a vertex u set $c(u) = \sigma(u) - d(u)$, and for a set of vertices U define $c(U) = \sum_{u \in U} c(u)$. So we can write the inequality above as $c(W) \leq \beta + \zeta |W|$.

In the following we use the fact that all a_e are large enough to find a bound for the sum $\sum_{e \in E(W)} a_e^{-1}$. To this aim, recall from (3) that $a_e = \left(\frac{3}{2}\beta + \frac{9}{2}\zeta\right) - c(u) - c(v)$ for all edges $e = uv$ in H . This gives

$$\sum_{e \in E(W)} a_e = \left(\frac{3}{2}\beta + \frac{9}{2}\zeta\right) |E(W)| - \sum_{u \in W} c(u) d_{H[W]}(u).$$

Since $d_{H[W]}(u) \leq d(u) = \sigma(u) - c(u) \leq \beta - c(u)$, we have

$$\sum_{e \in E(W)} a_e \geq \left(\frac{3}{2}\beta + \frac{9}{2}\zeta\right) |E(W)| - \beta c(W) + \sum_{u \in W} c(u)^2.$$

Set $q = \frac{3}{2}\beta + \frac{9}{2}\zeta$ and $p = \min_{uv \in E(W)} \{q - c(u) - c(v)\}$. This means that $q - p = \max_{uv \in E(W)} \{c(u) + c(v)\}$. Let $e = uv$ be an edge in $E(W)$ so that $c(u) + c(v) = q - p$. Then $c(u)^2 + c(v)^2 \geq \frac{1}{2}(q - p)^2$, and hence we can be sure that

$$\sum_{e \in E(W)} a_e \geq q |E(W)| - \beta c(W) + \frac{1}{2}(q - p)^2.$$

We now use this inequality and the following claim to bound $\sum_{e \in E(W)} a_e^{-1}$.

Claim 3.5. — Let r_1, \dots, r_m be m real numbers such that $0 < p \leq r_1, \dots, r_m \leq q$ and $\sum_{1 \leq i \leq m} r_i \geq qm - (q - p)S$, for some $S \geq 0$. Then we have $\sum_{1 \leq i \leq m} r_i^{-1} \leq \frac{S}{p} + \frac{m - S}{q}$.

Proof The result is trivial if $p = q$, so suppose $p < q$. For any $1 \leq i \leq m$, set $c_i = \frac{q - r_i}{q - p}$. Now we have $0 \leq c_i \leq 1$ for all $1 \leq i \leq m$, and $\sum_{1 \leq i \leq m} c_i \leq S$. Since the function $x \mapsto \frac{1}{x}$ is convex, we have that for $1 \leq i \leq m$,

$$\frac{1}{r_i} = \frac{1}{q - c_i(q - p)} = \frac{1}{c_i p + (1 - c_i)q} \leq c_i \frac{1}{p} + (1 - c_i) \frac{1}{q} = c_i \left(\frac{1}{p} - \frac{1}{q}\right) + \frac{1}{q}.$$

As a consequence,

$$\sum_{1 \leq i \leq m} \frac{1}{r_i} \leq \left(\frac{1}{p} - \frac{1}{q}\right) \sum_{1 \leq i \leq m} c_i + \frac{m}{q} \leq \left(\frac{1}{p} - \frac{1}{q}\right) S + \frac{m}{q} \leq \frac{S}{p} + \frac{m - S}{q}. \quad \square$$

We set $R = \beta c(W) - \frac{1}{2}(q - p)^2$ and $S = \frac{R}{q - p}$. Using Claim 3.5, at this point we have

$$\sum_{e \in E(W)} \frac{1}{a_e} \leq \frac{S}{p} + \frac{|E(W)| - S}{q} = \frac{S(q - p)}{pq} + \frac{|E(W)|}{q} = \frac{R}{pq} + \frac{2|E(W)|}{3\beta + 9\zeta}.$$

Notice that by condition (H3') of Lemma 2.12, $2|E(W)| \leq \sum_{u \in W} \sigma(u) - 2c(W) + \zeta|W| \leq \beta|W| - 2c(W) + \zeta|W|$. Hence we find

$$(4) \quad \sum_{e \in E(W)} \frac{1}{a_e} \leq \frac{\beta|W|}{3\beta + 9\zeta} + \frac{R}{pq} - \frac{2c(W)}{3\beta + 9\zeta} + \frac{\zeta|W|}{3\beta + 9\zeta}.$$

Claim 3.6. — We have $\frac{R}{pq} - \frac{2c(W)}{3\beta + 9\zeta} + \frac{\zeta|W|}{3\beta + 9\zeta} \leq \frac{\zeta}{\beta + 3\zeta}|W|$.

Proof. — Since $q = \frac{3}{2}\beta + \frac{9}{2}\zeta$, we only have to prove that $\frac{2R}{p} - 2c(W) \leq 2\zeta|W|$.

Let us write $q - p = \alpha\beta$, and so $p = \frac{1}{2}(3 - 2\alpha)\beta + \frac{9}{2}\zeta$ and $R = \beta c(W) - \frac{1}{2}\alpha^2\beta^2$. We have

$$\frac{2R}{p} - 2c(W) = \frac{2\beta c(W)}{p} - \frac{\alpha^2\beta^2}{p} - 2c(W).$$

If $p \geq \beta$, this expression is negative, so we can assume that $p < \beta$. In this case, using that $c(W) \leq \beta + \zeta|W|$, we have

$$\begin{aligned} \frac{2R}{p} - 2c(W) &= \frac{2\beta c(W)}{p} - \frac{\alpha^2\beta^2}{p} - 2c(W) \\ &= 2c(W) \frac{\beta - p}{p} - \frac{\alpha^2\beta^2}{p} \leq \frac{\beta}{p}(2\beta - 2p - \alpha^2\beta) + 2\zeta|W| \frac{\beta - p}{p}. \end{aligned}$$

As $2p = (3 - 2\alpha)\beta + 9\zeta$, we have $2\beta - 2p - \alpha^2\beta = (-1 + 2\alpha - \alpha^2)\beta - 9\zeta = -(\alpha - 1)^2\beta - 9\zeta < 0$. Since $p = a_e$ for some edge e , we have $p \geq \frac{1}{2}\beta$ by observation (c). Hence, $(\beta - p)/p \leq 1$ and we can conclude that $2R/p - 2c(W) \leq 2\zeta|W|$, which completes the proof of the claim. \square

Combining (4) and Claim 3.6, we obtain

$$\sum_{e \in E(W)} y_e = \sum_{e \in E(W)} \frac{1}{a_e} \leq \frac{\beta|W|}{3\beta + 9\zeta} + \frac{\zeta|W|}{\beta + 3\zeta} = \frac{\beta + 3\zeta}{3\beta + 9\zeta}|W| = \frac{1}{3}|W|.$$

Since $|W| \geq 3$, we have $\frac{1}{3}|W| \leq \frac{1}{2}(|W| - 1)$, which completes the proof of the lemma. \blacksquare

4. Proof of Theorem 1.3

We use the notation and terminology from Section 2.

We start similarly to the proof of Theorem 1.2 in Subsection 2.2. Suppose Theorem 1.3 is false. Then there exists a surface S such that for any β_S, γ_S we can find $\beta \geq \beta_S$ and a graph G , with a Σ -system of width at most β , such that $\omega(G; \Sigma) > \frac{3}{2}\beta + \gamma_S$. Let $\zeta_S^* = 132(3 - \chi(S))$ be as given in Lemma 2.4. We take $\zeta_S = \zeta_S^*$, $\beta_S = \frac{2}{3}(\zeta_S^*)^2 = 11616(3 - \chi(S))^2$, and $\gamma_S = \frac{1}{2}\zeta_S^* + 10 = 208 - 66\chi(S)$. Note that $\chi(S) \leq 2$, so $\beta_S \geq 11616$.

By assumption, there exist $\beta \geq \beta_S$ and a graph G , with a Σ -system of width at most β , containing a Σ -clique having more than $\frac{3}{2}\beta + \gamma_S$ vertices. Choose such graph G with the minimum number of vertices, and, with respect to that, with the maximum number of edges.

Similarly as in the proof of Theorem 1.2, we can assume G is connected, has at least 17424 vertices, and is edge-maximal with respect to being embeddable in S . By Lemma 2.3 we get that each vertex has degree at least three.

The following is an easy observation.

Claim 4.1. — For any vertex v , every Σ -clique containing v has size at most $1 + d^\Sigma(v)$.

Next we prove the following.

Claim 4.2. — Let adjacent vertices v, u satisfy $d(v) \leq 5$ and $d(u) \leq \zeta_S$. Then v is in every Σ -clique of size larger than $\frac{3}{2}\beta + \gamma_S$, and $d^\Sigma(v) \geq \frac{3}{2}\beta + \gamma_S$.

Proof. — The argument is similar to the one in Subsection 2.2.2: Construct a graph G_2 by contracting the edge vu into a new vertex w (removing multiple edges if they appear). Set $V_2 = (V \setminus \{v, u\}) \cup \{w\}$. Let $\Sigma_2(w) = \Sigma(u) \cup \Sigma(v) \setminus \{u, v\}$. For a vertex $t \in V_2 \setminus \{w\}$, if $\Sigma(t)$ contains u , then set $\Sigma_2(t) = (\Sigma(t) \setminus \{u, v\}) \cup \{w\}$; otherwise set $\Sigma_2(t) = \Sigma(t) \setminus \{v\}$. Note that G_2 is smaller than G and is still embeddable in S . Moreover, for every $t \in V_2 \setminus \{w\}$ we have $|\Sigma_2(t)| \leq |\Sigma(t)| \leq \beta$; while for w we have $|\Sigma_2(w)| \leq |\Sigma(u)| + |\Sigma(v)| \leq d_G(u) + d_G(v) \leq 5 + \zeta_S \leq \beta$.

By construction, it is easy to check that every Σ -clique in G not containing v corresponds to a Σ_2 -clique in G_2 of the same size. Since G was chosen as a smallest counterexample, this means that every Σ -clique in G of size larger than $\frac{3}{2}\beta + \gamma_S$ must contain v .

For the second part we use that G , as a counterexample, must contain Σ -cliques larger than $\frac{3}{2}\beta + \gamma_S$, whereas any Σ -clique in G containing v has size at most $1 + d^\Sigma(v)$. \square

We continue going through the cases of Lemma 2.4. If all vertices of G have degree at most ζ_S , then the number of Σ -neighbours of any vertex is at most $(\zeta_S)^2$. So the maximum size of a Σ -clique is at most $(\zeta_S)^2 + 1 \leq \frac{3}{2}\beta + 1 \leq \frac{3}{2}\beta + \gamma_S$, a contradiction.

Next suppose there is a vertex v of degree at most five with at most one neighbour of degree more than ζ_S . Then, since $\beta \geq \frac{2}{3}(\zeta_S^*)^2 \geq 8\zeta_S^*$, we have

$$d^\Sigma(v) \leq d(v) + \sum_{t \in N(v), v \in \Sigma(t)} (|\Sigma(t)| - 1) \leq 5 + 4(\zeta_S^* - 1) + (\beta - 1) = 4\zeta_S^* + \beta \leq \frac{3}{2}\beta.$$

But every vertex has degree at least three, hence v has a neighbour u of degree at most ζ_S^* . We obtain a contradiction with Claim 4.2.

Let X and Y be the two disjoint, non-empty, sets forming a very special ζ_S^* -pair in G satisfying (S3) in Lemma 2.4. For convenience, we repeat the essential properties of those sets:

- (i) *Every vertex in X has degree at least $\zeta_S^* + 1$. Every vertex $y \in Y$ has degree four, is adjacent to exactly two vertices of X , and the remaining neighbours of y have degree four as well.*
- (ii) *For all pairs of vertices $y, z \in Y$, if y and z are adjacent or have a common neighbour $w \notin X$, then $X^y = X^z$.*
- (iii) *For all non-empty subsets $W \subseteq X$, we have $e(W, Y \setminus Y^W) \leq e(W, Y \setminus Y^W) + \zeta_S^* |W|$.*

We can remove from X any vertex not adjacent to any vertex in Y .

We can use arguments similar to the first part of Subsection 2.2.3 to show the following.

Claim 4.3. — *For all $y \in Y$, we have that if $X^y = \{x_1, x_2\}$, then $y \in \Sigma(x_1) \cap \Sigma(x_2)$.*

Next, by (i), every $y \in Y$ has degree four and a neighbour u of degree four. From Claim 4.2 we can conclude:

Claim 4.4. — *For every $y \in Y$ we have that y is in every Σ -clique of size larger than $\frac{3}{2}\beta + \gamma_S$, and $d^\Sigma(y) \geq \frac{3}{2}\beta + \gamma_S$.*

Also by the properties of the vertices in Y according to (i) and (ii), we have for all $y \in Y$ and $X^y = \{x_1, x_2\}$,

$$\begin{aligned} d^\Sigma(y) &\leq 4 + 2 \cdot (4 - 1) + |\Sigma(x_1) \setminus \{y\}| + |\Sigma(x_2) \setminus \{y\}| - |Y^{\{x_1, x_2\}} \setminus \{y\}| \\ &= 9 + |\Sigma(x_1)| + |\Sigma(x_2)| - |Y^{\{x_1, x_2\}}| \end{aligned}$$

Here we use that by Claim 4.3 all vertices in $Y^{\{x_1, x_2\}}$ are contained in both $\Sigma(x_1)$ and $\Sigma(x_2)$; hence we can subtract the term $|Y^{\{x_1, x_2\}} \setminus \{y\}|$, since these vertices are counted twice in $|\Sigma(x_1) \setminus \{y\}| + |\Sigma(x_2) \setminus \{y\}|$. Since $|\Sigma(x_1)|, |\Sigma(x_2)| \leq \beta$, from Claim 4.4 we can conclude the following.

Claim 4.5. — *For every pair $x_1, x_2 \in X$ for which there is a $y \in Y$ with $X^y = \{x_1, x_2\}$, we have $|Y^{\{x_1, x_2\}}| \leq \frac{1}{2}\beta - \gamma_S + 9$.*

Since every vertex in Y is in every Σ -clique of size larger than $\frac{3}{2}\beta + \gamma_S$, and by the hypothesis there is at least one such clique, we must have that all pairs of vertices in Y are adjacent or appear together in some $\Sigma(v)$. By (ii), this proves that for every two vertices $y_1, y_2 \in Y$, we have $X^{y_1} \cap X^{y_2} \neq \emptyset$. As a consequence, if G_X denotes the graph with vertex set X in which two vertices are adjacent if they have a common neighbour in Y , then G_X is either a triangle or a star. (Here we use that we can assume all vertices in X to have at least one neighbour in Y .)

Case 1. G_X is a triangle.

Let $X = \{x_1, x_2, x_3\}$. This means that $Y = Y^{\{x_1, x_2\}} \cup Y^{\{x_1, x_3\}} \cup Y^{\{x_2, x_3\}}$, and so by Claim 4.5 we get $|Y| \leq \frac{3}{2}\beta - 3\gamma_S + 27$.

Since $Y^X = Y$ by definition of X , we have $e(X, Y \setminus Y^X) = 0$. So using the inequality in (iii) with $W = X$ leads to $e(X, V \setminus Y) \leq 3\zeta_S^*$. That means there must be x_{j_1} and x_{j_2} such that $e(\{x_{j_1}, x_{j_2}\}, V \setminus Y) \leq 2\zeta_S^*$. And so for $y \in Y^{\{x_{j_1}, x_{j_2}\}}$, we can estimate, using (i) and $|X| = 3$,

$$\begin{aligned} d^\Sigma(y) &\leq 2 + 2 \cdot (4 - 1) + |X| + (|Y| - 1) + e(\{x_{j_1}, x_{j_2}\}, V \setminus (X \cup Y)) \\ &\leq \frac{3}{2}\beta - 3\gamma_S + 37 + 2\zeta_S^*. \end{aligned}$$

But this contradicts Claim 4.4, since $4\gamma_S > 2\zeta_S^* + 37$.

Case 2. G_X is a star.

We denote by x the vertex of X corresponding to the centre of the star G_X , and by x_1, \dots, x_k , $k \geq 1$, the vertices of X corresponding to the leaves.

Using the inequality in (iii) with $W = X$ again, we get $e(X, V \setminus Y) \leq \zeta_S^* |X| = (k + 1)\zeta_S^*$. Since $X = \{x, x_1, \dots, x_k\}$, there must be an x_j such that $e(\{x_j\}, V \setminus Y) \leq \frac{1}{k}(k + 1)\zeta_S^* \leq 2\zeta_S^*$. Now for $y \in Y^{\{x, x_j\}}$, we can estimate

$$d^\Sigma(y) \leq 4 + 2 \cdot (4 - 1) + |(\Sigma(x) \cup \Sigma(x_j)) \setminus \{y\}| = 9 + |\Sigma(x)| + |\Sigma(x_j) \setminus \Sigma(x)|.$$

Since $Y \subseteq \Sigma(x)$, we have $|\Sigma(x_j) \setminus \Sigma(x)| \leq e(\{x_j\}, V \setminus Y) \leq 2\zeta_S^*$. Together with $|\Sigma(x)| \leq \beta$, this means $d^\Sigma(y) \leq \beta + 9 + 2\zeta_S^*$. This contradicts Claim 4.4, since $\frac{1}{2}\beta > 9 + 2\zeta_S^*$. ■

In the proof of Theorem 1.3, we used $\beta_S = 11616(3 - \chi(S))^2$ and $\gamma_S = 208 - 66\chi(S)$. Since the sphere \mathbb{S}^2 has $\chi(\mathbb{S}^2) = 2$, following the proof above means we can obtain $\beta_P = 11616$ and $\gamma_P = 76$ for the planar case. But it is clear that these values are far from best possible. Using more careful estimates in the proof above and more careful reasoning in certain parts of the proof of Lemma 2.4 can give significantly smaller values. Since our first goal is to show that we can obtain constant values for these results, we do not pursue this further.

5. Concluding Remarks and Discussion

5.1. About the Proof. — The proof of our main theorem for major parts follows the same lines as the proof of Theorem 1.6 in [10]. In particular, the proof of that theorem also starts with a structural lemma comparable to Lemma 2.4, uses the structure of the graph to reduce the problem to edge-colouring a specific multigraph, and then applies (and extends) Kahn’s approach to that multigraph. Of course, a difference is that Theorem 1.6 only deals with list colouring the square of a graph, but it is probably possible to generalise the whole proof to the case of list Σ -colouring. Nevertheless, there are some important differences in the proofs we feel deserve highlighting.

Lemma 2.4 is stronger than the comparable [10, Lemma 3.3]. We obtain a set Y of vertices with degree four and with a very specific structure of their neighbourhoods. This structure allows us to construct a multigraph H so that a standard list edge-colouring of H provides the information to colour the vertices in Y (see Lemma 2.7). In the lemma in [10], the vertices in the comparable set Y are only guaranteed to have degree at most $\Delta^{1/4}$, and knowledge about their neighbourhood is far sketchier. This means that the translation to list edge-colouring of a multigraph is not so clean; apart from the normal condition in the list edge-colouring of H (that adjacent edges need different colours), for each edge there may be up to $O(\Delta^{1/2})$ non-adjacent edges that also need to get a different colour. In particular this means that in [10], Kahn’s result in Theorem 2.11 cannot be used directly. Instead, a new, stronger, version has to be proved that can deal with a certain number of non-adjacent edges that need to be coloured differently. Lemma 2.4 allows us to use Kahn’s Theorem directly.

A second aspect in which our Lemma 2.4 is stronger is that in the final condition (S3), we have an ‘error term’ that is a constant times $|W|$. In [10] the comparable term is $\Delta^{9/10}|W|$, where Δ is the maximum degree of the graph. This in itself already means that the approach in [10] at best can give a bound of the type $\frac{3}{2}\Delta + o(\Delta)$. The fact that we cannot do better with the stronger structural result is because of the limitations of Kahn’s Theorem, Theorem 2.11. If it would be possible to replace the condition in that theorem by a condition of the form ‘the vector $\vec{x} = (x_e)$ with $x_e = \frac{1}{|L(e)| - K}$ for all $e \in E(H)$ is an element of $\mathcal{MP}(H)$ ’, where K is some positive constant, the work in this paper would directly give an improvement for the bound in Theorem 1.2 to $\frac{3}{2}\beta + O(1)$. Note that our version of Lemma 2.12 is also already strong enough to support that case.

Lemma 2.4 also allows us to prove a bound $\frac{3}{2}\beta + O(1)$ for the Σ -clique number in Theorem 1.3. The important corollary that the square of a graph embeddable on a fixed surface has clique number at most $\frac{3}{2}\Delta + O(1)$ would have been impossible without the improved bound in the lemma.

Also Lemma 2.12 is stronger than its compatriot [10, Lemma 5.9]. The lemma in [10] only deals with the case $d_G(v) = \beta$ for all vertices v in H . Because of this, it can only be applied to the case that all vertices in H have maximum degree $\Delta(G)$ in G . Some non-trivial trickery then has to be used to deal with the case that there are vertices in H of degree less than $\Delta(G)$ in G . Moreover, the proof of Lemma 2.12 is completely different from the proof in [10]. We feel that our new proof is more natural and intuitive, giving a clear relation between the lower bounds on the sizes of the lists and the upper bound of the sum of their inverses. The proof in [10] is more ad-hoc, using some non-obvious distinction in a number of different cases, depending on the size of W and the degrees of some vertices in W .

5.2. Further Work. — We feel that our work is just the beginning of the study of general Σ -colouring problems. It should be possible to obtain deeper results taking into account the structure of the Σ -system, and not just the sizes of the sets $\Sigma(v)$. The following easy result is an example of this.

Recall that a graph is *q-degenerate* if there exists an ordering v_1, v_2, \dots, v_n of the vertices such that every v_i has at most q neighbours in $\{v_1, \dots, v_{i-1}\}$. A class of graphs is degenerate if there is some q such that every graph in the class is q -degenerate. Examples of degenerate graph classes are graphs embeddable on a fixed surface, and proper minor-closed classes.

Proposition 5.1. — *For any degenerate graph class \mathcal{F} , there exists a constant $c_{\mathcal{F}}$ such that the following holds. Let G be a graph in \mathcal{F} , together with a Σ -system so that $\Sigma(u) \cap \Sigma(v) = \emptyset$ for every two distinct vertices u, v . Then $ch(G; \Sigma) \leq \Delta(G; \Sigma) + c_{\mathcal{F}}$.*

Proof Suppose every graph in \mathcal{F} is q -degenerate, and set $c_{\mathcal{F}} = q + 1$. For a graph G in \mathcal{F} , take an ordering v_1, \dots, v_n of its vertices such that each v_i has at most q neighbours in $\{v_1, \dots, v_{i-1}\}$. We greedily colour the vertices v_1, \dots, v_n in G in that order.

Note that by the hypothesis, each vertex v has at most one neighbour w with $v \in \Sigma(w)$. When colouring the vertex v_i , we need to take into account its neighbours in $\{v_1, \dots, v_{i-1}\}$, plus the vertices in $\Sigma(w) \cap \{v_1, \dots, v_{i-1}\}$ for a vertex w with $v_i \in \Sigma(w)$ (where that vertex w can be in $\{v_{i+1}, \dots, v_n\}$). By construction of the ordering, there are at most q neighbours of v_i in $\{v_1, \dots, v_{i-1}\}$. And a vertex w with $v_i \in \Sigma(w)$ has at most $|\Sigma(w)| \leq \Delta(G; \Sigma)$ vertices in $\Sigma(w) \cap \{v_1, \dots, v_{i-1}\}$. So the total number of forbidden colours when colouring v_i is at most $\Delta(G; \Sigma) + q$. Since each vertex has $\Delta(G; \Sigma) + q + 1$ colours available, the greedy algorithm will always find a free colour. ■

We think that it is possible to combine our main theorem and the theorem above in the following way. For a Σ -system for a graph G , let $k(G; \Sigma)$ be the maximum of $|\Sigma(u) \cap \Sigma(v)|$ over all pairs u, v of distinct vertices.

Conjecture 5.2. — *Let S be a fixed surface. Then there exists a constant c_S such that for all graphs G embeddable on S , with a Σ -system, we have*

$$ch(G; \Sigma) \leq \Delta(G; \Sigma) + k(G; \Sigma) + c_S.$$

This conjecture would fit with our current proof of Theorem 1.2, the main part of which is a reduction of the original problem to a list edge-colouring problem. For this approach, Shannon’s Theorem [31] that a multigraph with maximum degree Δ has an edge-colouring using at most $\lfloor \frac{3}{2} \Delta(G) \rfloor$ colours, forms a natural base for the bounds conjectured in Conjecture 1.1. If the relation between colouring the square of graphs embeddable on a fixed surface and edge-colouring multigraphs holds in a stronger sense, then Conjecture 5.2 forms a logical extension of Vizing’s Theorem [32] that a multigraph with maximum degree Δ and maximum edge-multiplicity μ has an edge-colouring with at most $\Delta + \mu$ colours.

In Borodin *et al.* [5], a weaker version of Conjecture 5.2 for cyclic colouring of plane graphs was proved. Recall that if G^P is a plane graph, then Δ^* is the maximum number of vertices in a face. Let k^* denote the maximum number of vertices that two faces of G^P have in common.

Theorem 5.3 (Borodin, Broersma, Glebov & Van den Heuvel [5])

For a plane graph G^P with $\Delta^ \geq 4$ and $k^* \geq 4$ we have $\chi^*(G^P) \leq \Delta^* + 3k^* + 2$.*

Σ -Colouring and Minor-Closed Classes. — It seems natural to expect that our work on graphs embeddable in a fixed surface can be extended to arbitrary proper minor-closed classes of graphs. Compare our main Theorem 1.2 with Theorem 1.6, the main result from [10]. But there exist some obstacles to a direct generalisation.

It is easy to show that if a graph G is q -degenerate, then its square is $((2q - 1)\Delta(G))$ -degenerate. It is well-known, see e.g. [22], that for every proper minor-closed family \mathcal{F} , there is a constant $C_{\mathcal{F}}$ such that every graph in \mathcal{F} is $C_{\mathcal{F}}$ -degenerate. Hence G^2 is $((2C_{\mathcal{F}} - 1)\Delta(G))$ -degenerate, and so for every $G \in \mathcal{F}$ we have $ch(G^2) \leq (2C_{\mathcal{F}} - 1)\Delta(G) + 1$.

For Σ -colouring, there is no comparable upper bound on $ch(G; \Sigma)$ in terms of the degeneracy of G and $\Delta(G; \Sigma)$. To see this, let G be the graph obtained from the complete graph K_n , $n \geq 4$, by subdividing all edges of K_n once. For a vertex v corresponding to an original vertex in K_n , set $\Sigma(v) = \emptyset$; while for a “new” vertex v of degree two, set $\Sigma(v) = N_G(v)$. Then we have that G is 2-degenerate and $\Delta(G; \Sigma) = 2$, but $ch(G; \Sigma) = n$.

Nevertheless, combining the Robertson and Seymour graph minor structure theorem [28] with our main theorem on graphs embeddable in bounded genus surfaces, one can fairly easily obtain the following.

Theorem 5.4. — *Let \mathcal{F} be a proper minor-closed family of graphs. Then there exist constants $C_{\mathcal{F}}$ and $c_{\mathcal{F}}$ such that the following holds: For any graph G in \mathcal{F} with a Σ -system, we have $ch(G; \Sigma) \leq C_{\mathcal{F}}\Delta(G; \Sigma) + c_{\mathcal{F}}$.*

Giving more details of the ideas of the proof would require a number of additional definitions, and is beyond the scope of this short discussion. It would be interesting to find a proof of this theorem that does not require the full force of the graph minor structure theorem.

Also finding the smallest possible constant $C_{\mathcal{F}}$ for certain minor-closed families \mathcal{F} appears an interesting question. Theorem 1.6 clearly suggests that if \mathcal{F}_k denotes the class of $K_{3,k}$ -minor free graphs ($k \geq 3$), $C_{\mathcal{F}_k}$ should be equal to $3/2$.

6. Kahn’s Work on List Edge-Colourings

As mentioned earlier, Theorem 2.11 is not explicitly stated in [17], but is implicit in the proof of the main result of that paper. In this final section, we give an overview of how this theorem can be obtained from the ideas in Kahn’s paper.

The main result in [17] is that the list chromatic index is asymptotically equal to the fractional chromatic index of a multigraph.

Theorem 6.1 (Kahn [17]). — *For all $\varepsilon > 0$, there exists a Δ_{ε} such that for all $\Delta \geq \Delta_{\varepsilon}$ the following holds. If H is a multigraph with maximum degree at most Δ , then*

$$\chi'_f(H) \leq \chi'(H) \leq ch'(H) \leq (1 + \varepsilon)\chi'_f(H).$$

Here $\chi'(H)$ is the normal chromatic index (or edge-chromatic number) of H , $\chi'_f(H)$ is the fractional chromatic index of H , and $ch'(H)$ is the list chromatic index of H . The crucial step to relate this result to the matching polytope $\mathcal{MP}(H)$ is the following well-known characterisation of the fractional chromatic index:

$$\chi'_f(H) = \min\{\gamma > 0 \mid \text{the vector } (x_e)_{e \in E(H)} \text{ with } x_e = \gamma^{-1} \text{ is in } \mathcal{MP}(H)\}.$$

So Theorem 6.1 is just a special case of Theorem 2.11 if we set $|L(e)| \geq \frac{\chi'_f(H)}{1-\delta}$ for all edges e . (The second condition of Theorem 2.11 is automatically satisfied in that case, since trivially $\chi'_f(H) \geq \Delta(H)$.)

In order to prove Theorem 6.1, Kahn describes a randomised iterative procedure that colours the edges of H in a number of stages. During this procedure, the lists of available colours for each edge will change, and the lists will not be the same size for the uncoloured edges. This is why, roughly speaking, Kahn’s actual proof deals with the more general case, as described in Theorem 2.11.

In order to give the reader a better understanding of the background of Kahn’s approach, we give an overview of the crucial elements in the following subsections.

6.1. Hardcore Distributions. — Hardcore distributions are distributions that originally arose in Statistical Physics, and that satisfy very natural conditions and generally provide strong independence properties allowing good sampling from a given family. Given a family of subsets \mathcal{F} of a given set \mathcal{E} , a natural way of picking at random an element of \mathcal{F} (or, in other words, a probability distribution on \mathcal{F}) is as follows.

Let us suppose that each element e of \mathcal{E} has been assigned a positive weight λ_e . Then we pick each element $M \in \mathcal{F}$ with probability proportional to $\prod_{e \in M} \lambda_e$. More precisely, the probability P_M of picking $M \in \mathcal{F}$ at random is given by

$$P_M = \frac{\prod_{e \in M} \lambda_e}{\sum_{M' \in \mathcal{F}} \prod_{e \in M'} \lambda_e}.$$

We define the vector $\vec{x} = (x_e)_{e \in \mathcal{E}}$ by setting $x_e = \sum_{M \in \mathcal{F}, e \in M} P_M$. It is clear that x_e is the probability that a given random element of \mathcal{F} contains the element e . The probability distribution $\{P_M\}$ is called a *hardcore* distribution with *activities* $\{\lambda_e\}$ and *marginals* $\{x_e\}$. The vector \vec{x} is called the *marginal vector* associated with the hardcore distribution $\{P_M\}$.

Given a vector \vec{x} , it is not always true that \vec{x} is the marginal vector of some hardcore distribution. Indeed if $\mathcal{P}(\mathcal{F})$ denotes the polytope defined by taking the convex hull of the characteristic vectors of the elements of $\mathcal{F}^{(1)}$, then the marginal vector \vec{x} of a hardcore distribution is in $\mathcal{P}(\mathcal{F})$:

$$\vec{x} = \sum_{M \in \mathcal{F}} P_M \mathbf{1}_M.$$

This provides a necessary condition for a vector to be the marginal vector of a hardcore distribution. It is not difficult to prove that the activities λ_e corresponding to \vec{x} , if they exist, are unique.

From now on, let H be a given multigraph. We recall that $\mathcal{M}(H)$ and $\mathcal{MP}(H)$ are the family of matchings and the matching polytope of H , respectively. (So $\mathcal{M}(H)$ will play the role of the family \mathcal{F} from above. And using the notation from above means $\mathcal{MP}(H) = \mathcal{P}(\mathcal{M}(H))$.)

We have the following theorem relating the matching polytope and hardcore distributions.

⁽¹⁾Recall that the characteristic vector, $\mathbf{1}_M$, of a given element $M \in \mathcal{F}$ is the $|\mathcal{E}|$ -dimensional vector $(y_e)_{e \in \mathcal{E}}$ such that $y_e = 1$ if $e \in M$ and $y_e = 0$ otherwise.

Theorem 6.2 (Lee [20], Rabinovich *et al.* [27]). — For a given real number $0 < \delta < 1$, suppose \vec{x} is a vector in $(1 - \delta)\mathcal{MP}(H)$, for some multigraph H . Then there exists a unique family of activities λ_e such that \vec{x} is the marginal vector of the hardcore distribution defined by the λ 's. The hardcore distribution $\{P_M\}_{M \in \mathcal{M}(H)}$ is the unique distribution maximising the entropy function

$$\mathcal{H}(Q_M) = - \sum_{M \in \mathcal{M}(H)} Q_M \log(Q_M)$$

among all the distributions $\{Q_M\}_{M \in \mathcal{M}(H)}$ satisfying $\vec{x} = \sum_{M \in \mathcal{M}(H)} Q_M \mathbf{1}_M$.

Kahn and Kayll proved in [18] a family of results, resulting in a long-range independence property for the hardcore distributions defined by a marginal vector \vec{x} inside $(1 - \delta)\mathcal{MP}(H)$, see [17]. We refer to the original papers of Kahn [16, 17] and Kahn and Kayll [18], and the book by Molloy and Reed [24] for more on these issues. We settle here for citing the following lemma.

Lemma 6.3 ([18, Lemma 4.1]). — For every δ , $0 < \delta < 1$, there is a $\rho_\delta > 0$ such that if $\{P_M\}$ is a hardcore distribution with marginal vector $\vec{x} \in (1 - \delta)\mathcal{MP}(H)$, then for all $u, v \in V(H)$,

$$\Pr(M \text{ does not touch } u \text{ and } v) > \rho_\delta.$$

6.2. Hardcore Distributions and Edge-Colouring. — We present here Kahn's algorithm for list edge-colouring of multigraphs first introduced and analysed in [17]. We continue to use the notation of the previous subsection. In particular, we suppose that H is a multigraph and L a list assignment of colours to the edges of H so that the conditions of Theorem 2.11 are satisfied. By Lemma 6.2 there exists a hardcore distribution $\{P_M\}$ with marginals $\{|L(e)|^{-1}\}_{e \in E(H)}$, which in addition satisfies the property of Lemma 6.3. Let $\{\lambda_e\}$ be the activities on the edges (which are unique by Theorem 6.2) corresponding to this distribution. An extra condition is indeed true: For every subgraph H^* of H it is possible to find a hardcore distribution $\{P_M^*\}$ with corresponding marginals $|L(e)|^{-1}$ for $e \in E(H^*)$. The corresponding activities λ_e^* will in general be different from the λ_e 's.

The algorithm works as follows: Let $\mathcal{L} = \bigcup_{e \in E(H)} L(e)$ be the union of the colours in the lists.

For each colour α , let us define the *colour graph* H_α to be the graph containing all the edges whose lists contain the colour α . And denote by $\{\lambda_{\alpha,e}\}$ the activities producing the hardcore distribution with marginals $|L(e)|^{-1}$ for $e \in E(H_\alpha)$. The colouring procedure consists in a finite number of iterations of a procedure that we may call *naive colouring*. At step i of the iteration, we are left with subgraphs H_α^i containing some uncoloured edges whose lists contain the colour α . Of course we have $H_\alpha^i \subseteq H_\alpha^{i-1} \subseteq \dots \subseteq H_\alpha^0 = H_\alpha$.

The naive colouring procedure at step $i + 1$ consists of the following sub-steps.

- (a) For each colour $\alpha \in \mathcal{L}$, choose independently of the other colours a random matching $M_\alpha^{i+1} \subseteq E(H_\alpha^i)$. The distribution of the matchings is the hardcore distribution defined by the activities $\lambda_{\alpha,e}$ on the edges $e \in E(H_\alpha^i)$.
- (b) If an edge e is in one or more of the M_α^{i+1} 's, then choose one of the colours from those, chosen uniformly at random, and colour e with that colour.

- (c) For each colour α , form H_α^{i+1} by removing from H_α^i all the edges that received some colour at this stage, and all vertices that are incident to one of the edges coloured with α . (While removing a vertex, all the edges incident to it are of course removed as well.)

Note that the process above can be described also in terms of subgraphs H^i of the original multigraph H , where the edges of H^i are the edges that are still uncoloured after step i , and each edge e in H^i has a list of colours $L^i(e)$ formed by all colours α for which $e \in E(H_\alpha^i)$. Also note that the activities $\lambda_{\alpha,e}$ remain unchanged all through the process (but the edge sets on which they are applied change).

A sufficient number of iterations of the naive colouring procedure results in a graph H^I , consisting of all the uncoloured edges at this step, such that H^I has maximum degree T , for some integer T , and that the list sizes are at least $2T$, i.e., each uncoloured edge is in at least $2T$ of the H_α^I 's. (Remember that the conditions of Theorem 2.11 imply that the lists are quite large at the beginning.) At this stage it is easy to finish the procedure by a simple greedy algorithm.

The heart of the analysis of the above algorithm in Kahn's approach is the following strong lemma, the proof of which can be found in [17].

Lemma 6.4 (Kahn [17, Lemma 3.1]). — *For each $K > 0$ and $0 < \eta < 1$, there are constants $0 < \xi_{K,\eta} \leq \eta$ and $\Delta_{K,\eta}$ such that the following holds for all $\Delta \geq \Delta_{K,\eta}$. Let H be a multigraph with lists $L(e)$ of colours for each edge e . For each colour α , define the colour graph H_α as above. Finally, for each colour α we are given a hardcore distribution with activities $\{\lambda_{\alpha,e}\}_{e \in E(H_\alpha)}$ and marginals $\{x_{\alpha,e}\}_{e \in E(H_\alpha)}$. Suppose the following conditions are satisfied:*

- for every vertex v , $d_H(v) \leq \Delta$;
- for every colour α and edge $e \in E(H_\alpha)$, $\lambda_{\alpha,e} \leq \frac{K}{\Delta}$; and
- for every edge e , $1 - \xi_{K,\eta} \leq \sum_{\alpha \text{ in } L(e)} x_{\alpha,e} \leq 1 + \xi_{K,\eta}$.

Then with positive probability the naive colouring procedure described above gives matchings $M_\alpha \subseteq E(H_\alpha)$ for all colours α , so that if we set $H^ = H - \bigcup_{\alpha'} M_{\alpha'}$, $H_\alpha^* = H_\alpha - V(M_\alpha) - \bigcup_{\alpha'} M_{\alpha'}$, and form lists $L^*(e)$ for all edges $e \in E(H^*)$ by removing no longer allowed colours from $L(e)$, we have:*

- for every vertex v , $d_{H^*}(v) \leq \frac{1+\eta}{1+\xi_{K,\eta}} e^{-1} \Delta$; ⁽²⁾ and
- for every edge e in H^* , $1 - \eta \leq \sum_{\alpha \in L^*(e)} x_{\alpha,e}^* \leq 1 + \eta$.

Here $\{x_{\alpha,e}^*\}_{e \in E(H_\alpha^*)}$ are the marginals associated to $\lambda_{\alpha,e}$ in H_α^* .

In other words, the lemma guarantees that after one iteration of the naive colouring procedure, with positive probability the multigraph formed by the uncoloured edges has maximum degrees bounded by $\frac{1+\eta}{1+\xi_{K,\eta}} e^{-1} \Delta$, while the sum of the marginal probabilities $x_{\alpha,e}^*$ for every edge e will be close to 1.

⁽²⁾To avoid confusion between an edge 'e' and the base of the natural logarithms 2.718..., we will use the roman letter 'e' for the latter one.

In the next subsection we will combine all the strands and use the lemma above to conclude the proof of Theorem 2.11.

6.3. Completing the Proof of Theorem 2.11 — after Kahn. — Let $0 < \delta < 1$ and $\nu > 0$. Then we should prove the existence of a $\Delta_{\delta,\nu}$ such that for $\Delta \geq \Delta_{\delta,\nu}$ the following holds. Let H be a multigraph and L a list assignment of colours to the edges of H so that

- for every vertex v , $d_H(v) \leq \Delta$;
- for all edges $e \in E(H)$, $|L(e)| \geq \nu \Delta$;
- the vector $\vec{x} = (x_e)$ with $x_e = \frac{1}{|L(e)|}$ for all $e \in E(H)$ is an element of $(1 - \delta) \mathcal{MP}(H)$.

Then there should exist a proper edge-colouring of H , where each edge receives a colour from its own list.

For each colour α , define the colour graph H_α as in the previous subsection. For each colour α and edge e , set $x_{\alpha,e} = x_e = \frac{1}{|L(e)|}$, and let $\{\lambda_{\alpha,e}\}_{e \in E(H_\alpha)}$ be the activities associated with the marginals $x_{\alpha,e}$ on H_α .

Since for every edge e we have $\sum_{\alpha \in L(e)} x_{\alpha,e} = \sum_{\alpha \in L(e)} |L(e)|^{-1} = 1$, we certainly know that

- for every edge e and $\xi > 0$, $1 - \xi \leq \sum_{\alpha \in L(e)} x_{\alpha,e} \leq 1 + \xi$.

We next bound the activities $\lambda_{\alpha,e}$, using Lemma 6.3. First observe that for all α the vector $(x_{\alpha,e})_{e \in E(H_\alpha)}$ is in $(1 - \delta) \mathcal{MP}(H_\alpha)$. So by Lemma 6.3 there is a constant ρ_δ such that if M_α is chosen according to the hardcore distribution with marginals $\{x_{\alpha,e}\}$ on H_α , then for all $u, v \in E(H_\alpha)$ we have $\Pr(M_\alpha \text{ does not touch } u \text{ and } v) > \rho_\delta$. Let $e = uv$ be an edge of H_α . Then we have

$$x_{\alpha,e} = \Pr(M_\alpha \text{ contains } e) = \lambda_{\alpha,e} \cdot \Pr(M_\alpha \text{ does not touch } u \text{ and } v) > \lambda_{\alpha,e} \cdot \rho_\delta.$$

Given the fact that $x_{\alpha,e} = \frac{1}{|L(e)|}$ and $|L(e)| \geq \nu \Delta$, and setting $K_{\delta,\nu} = \frac{1}{\rho_\delta \nu}$, we infer that $\lambda_{\alpha,e} < \frac{x_{\alpha,e}}{\rho_\delta} \leq \frac{1}{\rho_\delta \nu \Delta} = \frac{K_{\delta,\nu}}{\Delta}$. We have shown that there exists a $K_{\delta,\nu} > 0$ such that

- for every colour α and edge $e \in E(H_\alpha)$, $\lambda_{\alpha,e} \leq \frac{K_{\delta,\nu}}{\Delta}$.

Suppose we repeat the naive colouring procedure from the previous subsection $s = s_{K_{\delta,\nu}}$ times (where $s_{K_{\delta,\nu}}$ is a fixed constant to be made more precise later). Let H^i be the subgraph of H formed by the edges that are as yet uncoloured at step i , and for each $e \in E(H^i)$ let $L^i(e)$ be the list of colours from $L(e)$ that are still allowed for e at that stage.

Let $\eta_s = 1 - e^{-1}$, and recursively in the up-to-down order for $i = s - 1, \dots, 1$, set $\eta_i = \xi_{K_{\delta,\nu}, \eta_{i+1}}$, where $\xi_{K_{\delta,\nu}, \eta_{i+1}}$ is the function given by Lemma 6.4. Let $\Delta_{\delta,\nu} = \max_{i=1, \dots, s} \Delta_{K_{\delta,\nu}, \eta_i}$ ($\Delta_{K_{\delta,\nu}, \eta_i}$ according to Lemma 6.4 again), and $\eta_0 = 0$. By applying Lemma 6.4 and the observations above, we can ensure inductively, starting from $i = 0$, that for $\Delta \geq \Delta_{\delta,\nu}$, with positive probability the following conditions are satisfied for all $i = 0, \dots, s$.

- For all vertices v , $d_{H^i}(v) \leq T_i$, where $T_0 = \Delta$ and $T_i = \frac{1 + \eta_i}{1 + \eta_{i-1}} e^{-1} T_{i-1}$ for $i \geq 1$; and

- For all edges $e \in E(H^i)$, $1 - \eta_i \leq \sum_{\alpha \in L^i(e)} x_{\alpha,e}^i \leq 1 + \eta_i$, where $\{x_{\alpha,e}^i\}_{e \in E(H^i)}$ are the marginals associated to the hardcore distribution with activities $\lambda_{\alpha,e}$ in H_α^i .

It follows that that after s steps, with positive probability we have

- for all vertices v , $d_{H^s}(v) \leq (2 - e^{-1})e^{-s} \Delta$; and
- for all edges $e \in E(H^s)$, $e^{-1} \leq \sum_{\alpha \in L^s(e)} x_{\alpha,e}^s \leq 2 - e^{-1}$.

We note that for an edge $e = uv$,

$$x_{\alpha,e}^s = \lambda_{\alpha,e} \cdot \mathbf{Pr}(M_\alpha^s \text{ does not touch } u \text{ and } v) \leq \lambda_{\alpha,e},$$

which implies that $x_{\alpha,e}^s \leq \lambda_{\alpha,e} \leq \frac{K_{\delta,\nu}}{\Delta}$. We infer that for all $e \in E(H^s)$,

$$|L^s(e)| = |\{\alpha \mid e \in E(H_\alpha^s)\}| \geq \frac{\Delta}{e K_{\delta,\nu}}.$$

Let $T = \frac{\Delta}{2e K_{\delta,\nu}}$. It is now clear that if we choose the value of s in such a way that $2e^{-s} \leq \frac{1}{2e K_{\delta,\nu}}$ (in other words, by setting $s = s_{K_{\delta,\nu}} \geq \ln(4K_{\delta,\nu}) + 1$), we can ensure with positive probability that

$$d_{H^s}(v) \leq T \quad \text{for all } v \in V(H^s) \quad \text{and} \quad |L^s(e)| \geq 2T \quad \text{for each } e \in E(H^s).$$

This finally shows that we can proceed using the greedy algorithm in H^s , in order to extend the resulting colouring from the naive colouring procedure in Subsection 6.2 to a colouring of the whole graph.

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